

# Edge-Aware Inter-domain Routing (EIR)

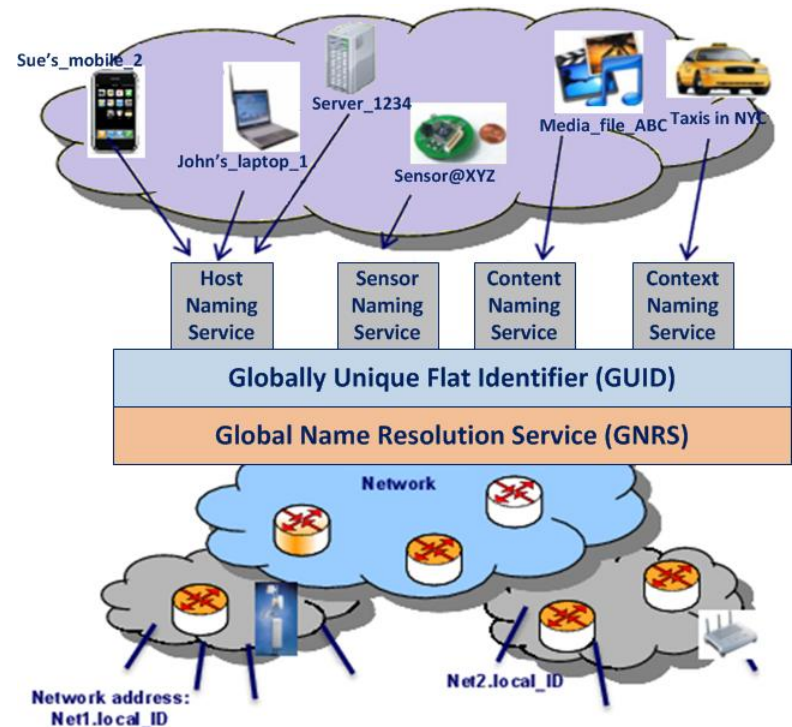
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# MobilityFirst Overview

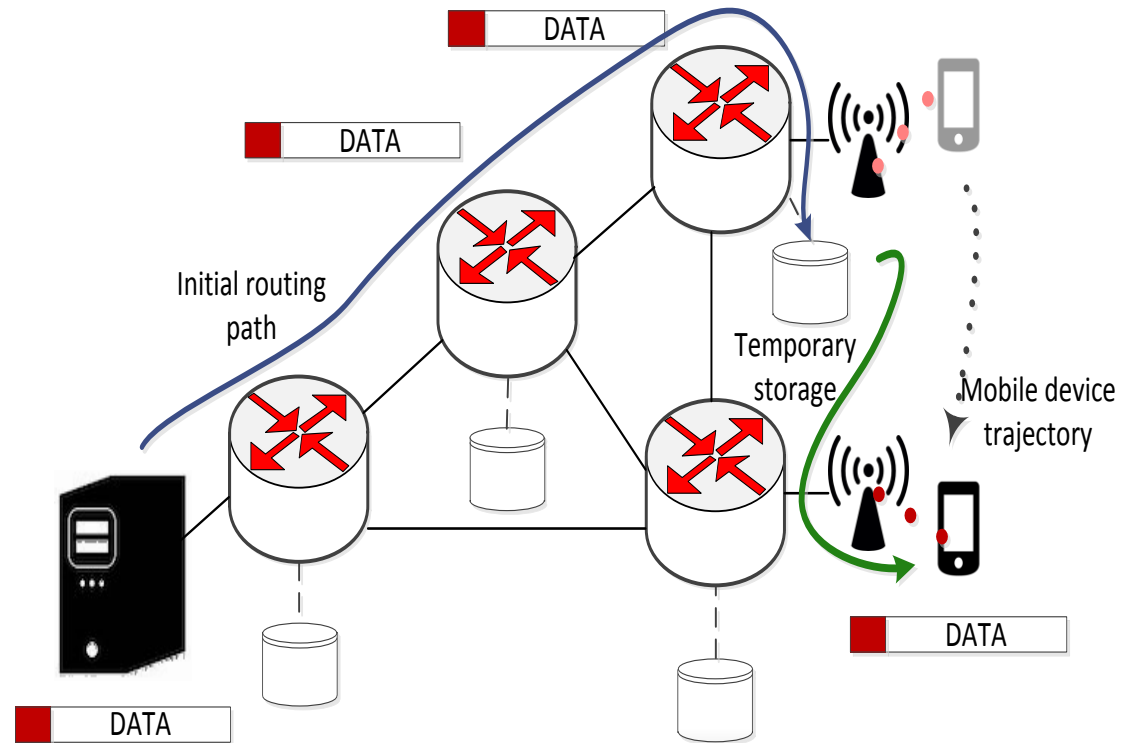
- Name based communication
  - Flat self-certifying globally unique identifier (GUID)
- Dynamic name resolution
  - Mapping of GUID to network address maintained at the global name resolution service (GNRS)
- In-network storage
- Hop by hop reliable transport



# Intra-domain routing in MF

Generalized storage aware routing (GSTAR):

- Proactive link state routing
- Link quality information in terms of expected time of transmission (ETT)
- Directly accessible in-network storage
- Store or forward decisions based on short term and long term ETTs



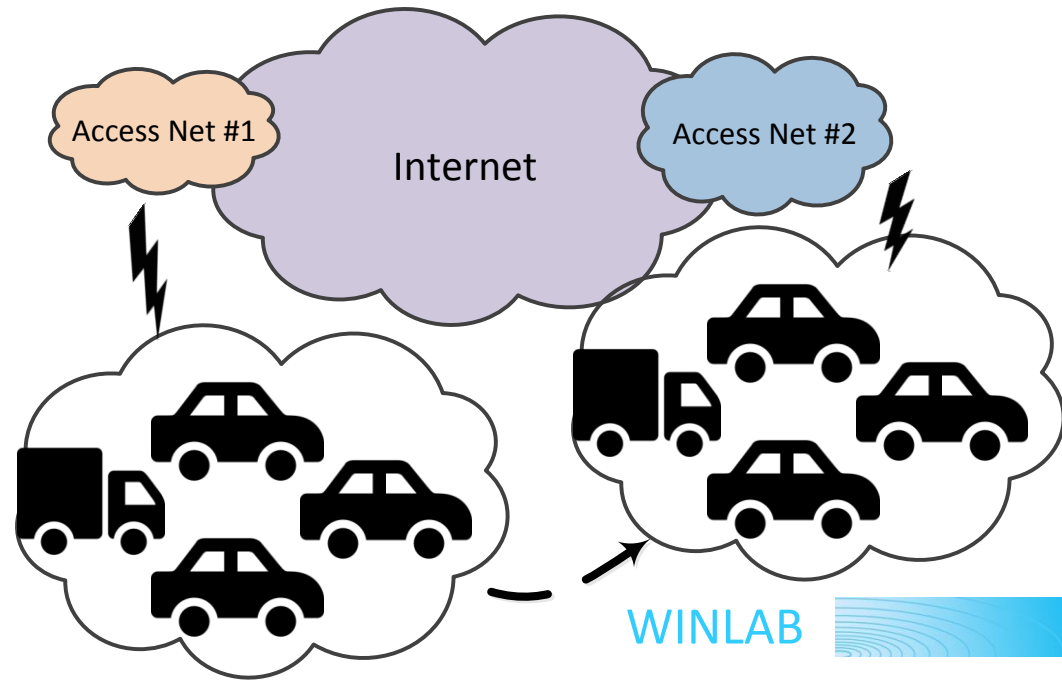
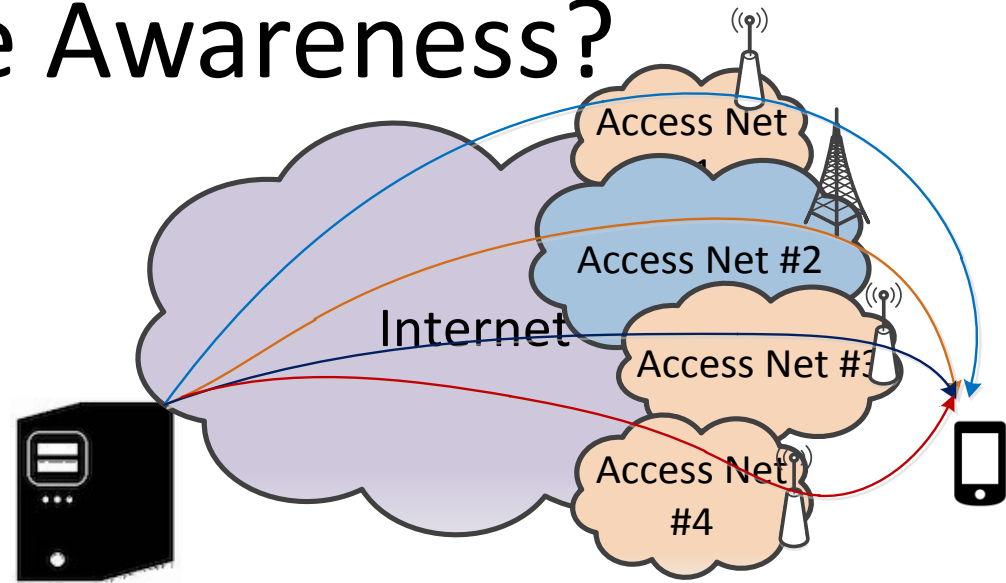
# Clean-slate Inter-domain Routing

- Current inter-domain routing has poor support for
  - Edge mobility
  - Dynamic AS formation
  - Multipath routing
  - Intra-domain network visibility
- EIR is a clean slate inter-domain routing framework designed for the MobilityFirst architecture



# Why Edge Awareness?

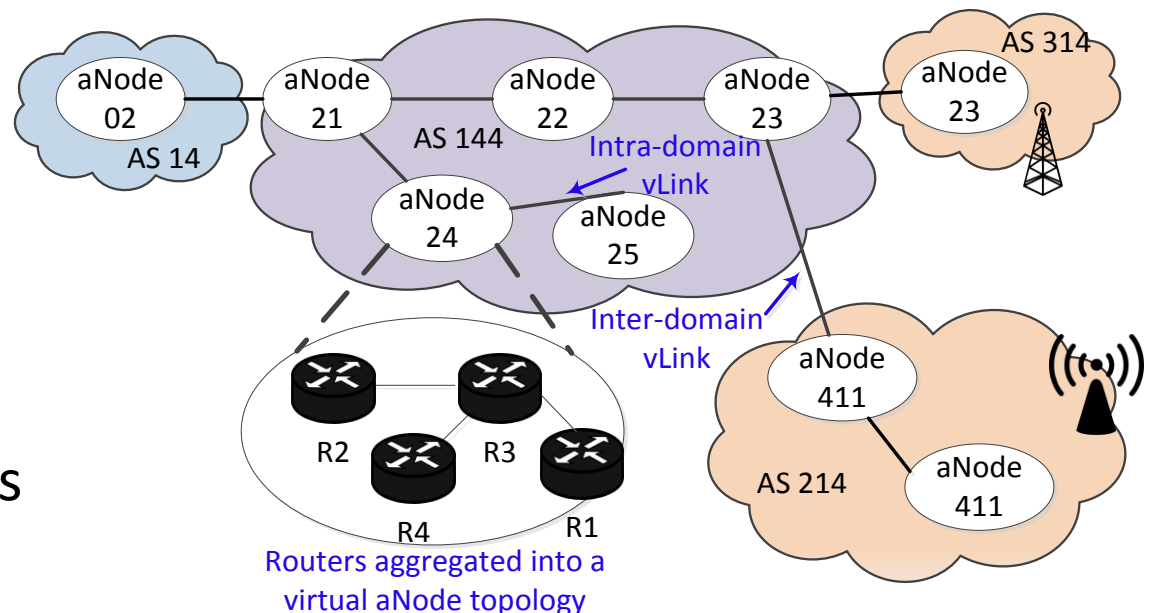
- Basic property of a mobile device:
  - Multi-radio capability
  - Multiple available paths at the edge
- Ad-hoc network formation without connectivity to the core
- Dynamic wireless edge peering with mobility



# EIR Design Features

## 1. Aggregation nodes (aNodes) and virtual links (vLinks)

- aNodes and vLinks as abstractions to expose internal network topology of an AS
  - Aggregated information of a collection of routers and links
  - Flexible aggregation granularity

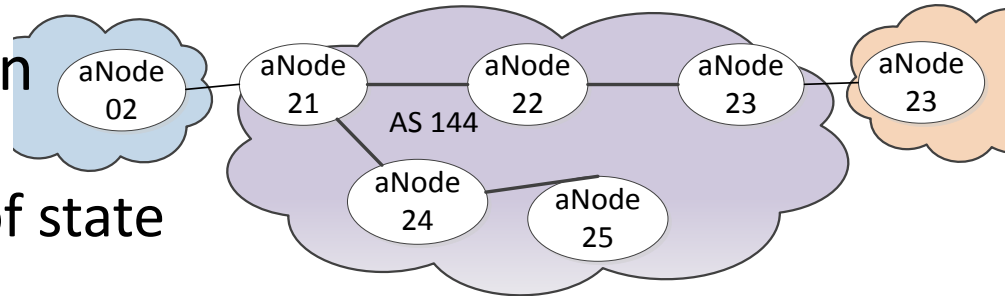


Aggregation of intra-network topology



# EIR Design Features

## 2. Network state dissemination



Aggregated view of a network

- Network-wide propagation of state packets (nSPs)
- nSPs contain
  - aNode and vLink connectivity info
  - vLink properties (Mbps, % avail, estimated time of transmission (ETT), functional parameters, etc.)
- Shortest path routing based on global aggregated topology

Msg_Type	AS_Num:Source_aNode	Hop_to_Src
Internal Topologies:		
aNode#1-vLink<B,V,A,L>-aNode#2		
aNode#2-vLink<B,V,A,L>-aNode#3		
...		
aNode#x-vLink<B,V,A,L>-aNode#y		
Neighbor Info:		
Neighbor_aNode#1-vLink<B,V,A,L>		
Neighbor_aNode#2-vLink<B,V,A,L>		
...		
Neighbor_aNode#z-vLink<B,V,A,L>		

Network state packet (nSP)



# EIR Design Features

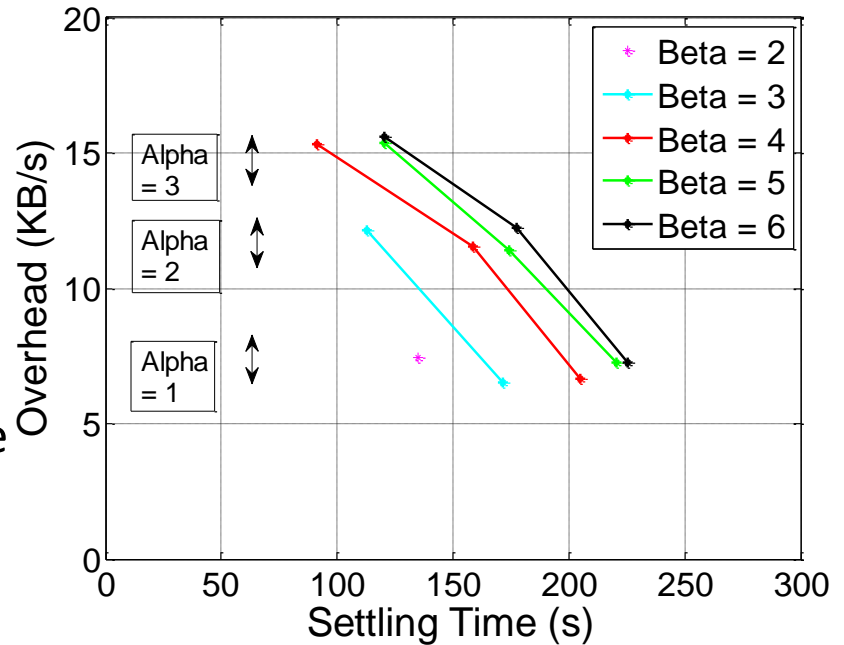
## 3. Telescopic flooding

- Border routers forward nSPs originated from other ASes in a telescopic manner
- The relaying rate is determined by the distance of the packet-source AS to the forwarding AS
- Nearby ASes will have a better view of the network than far-away ASes

$$y_1 = A$$

$$y_2 = A * \exp(x-1)$$

$$y_3 = \begin{cases} A, x < \alpha \\ A * \exp(x-\alpha), \alpha \leq x < \beta \\ A * \exp(\beta-\alpha), x \geq \beta \end{cases}$$



Parametric evaluation of overhead for const-exp-const function

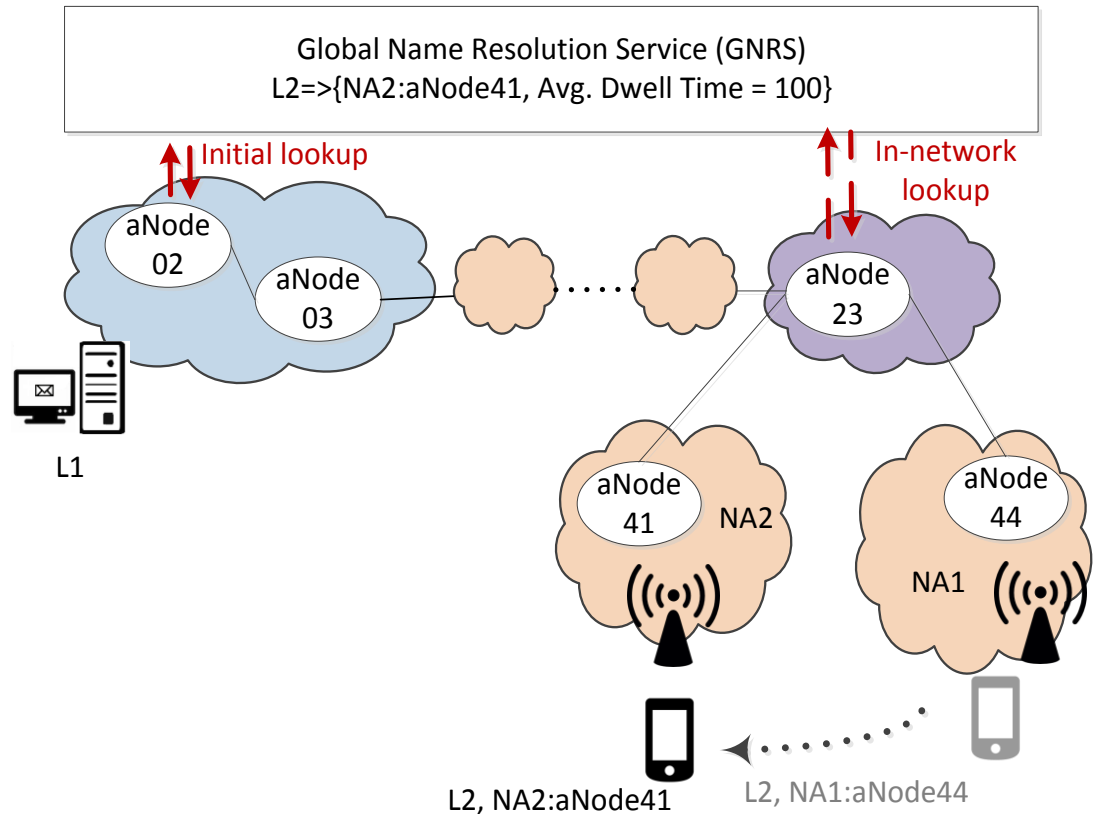




# EIR Design Features

## 4. Late binding support

- At source early binding with re-resolution on failure
- In-network binding of names (GUIDs) to locations (aNodes) closer to the edge for more up-to-date information

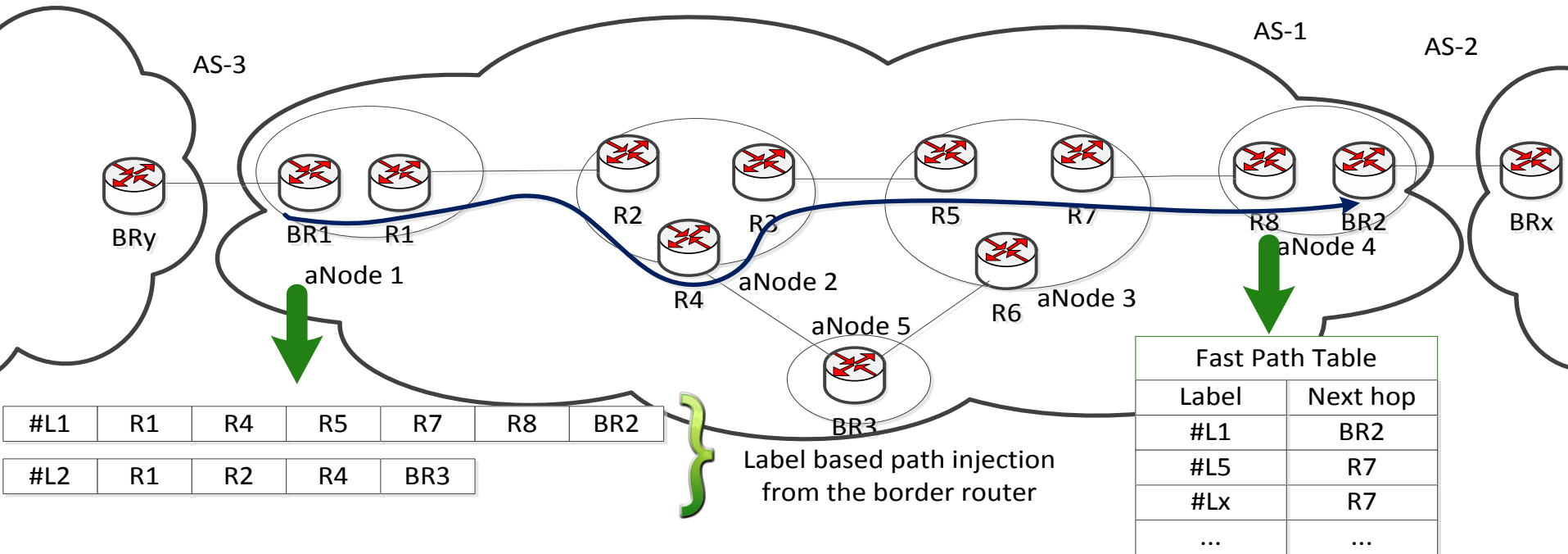


Mobility support through late binding



# EIR Design Features

## 5. Fast path for transit traffic

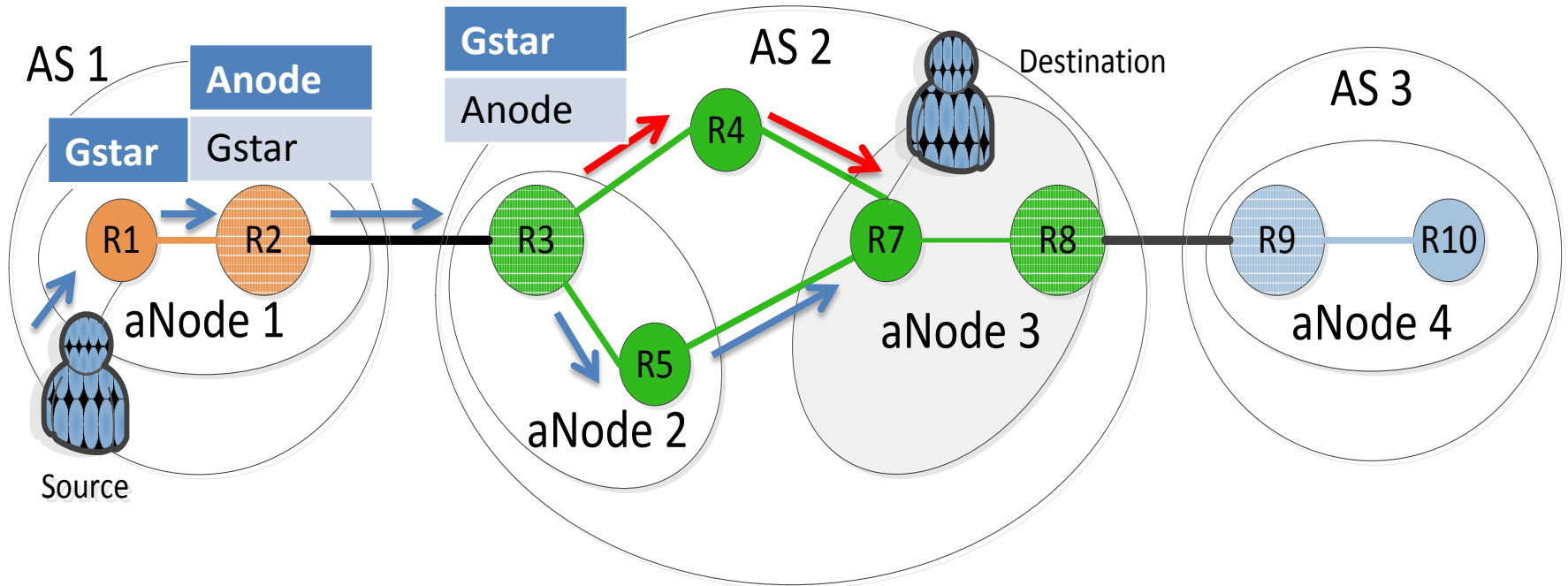


- Border routers determine cut-through paths based on transit and local policies
- Route injection packets are pushed by border routers to inject label info in fast path forwarding table
- Transit traffic at the ingress border router is 'marked' transit and appended with the label
- No inter-domain routing/policy info required at the internal routers



# EIR Routing Scenarios

## 1. Source in AS 1 and destination in AS 2

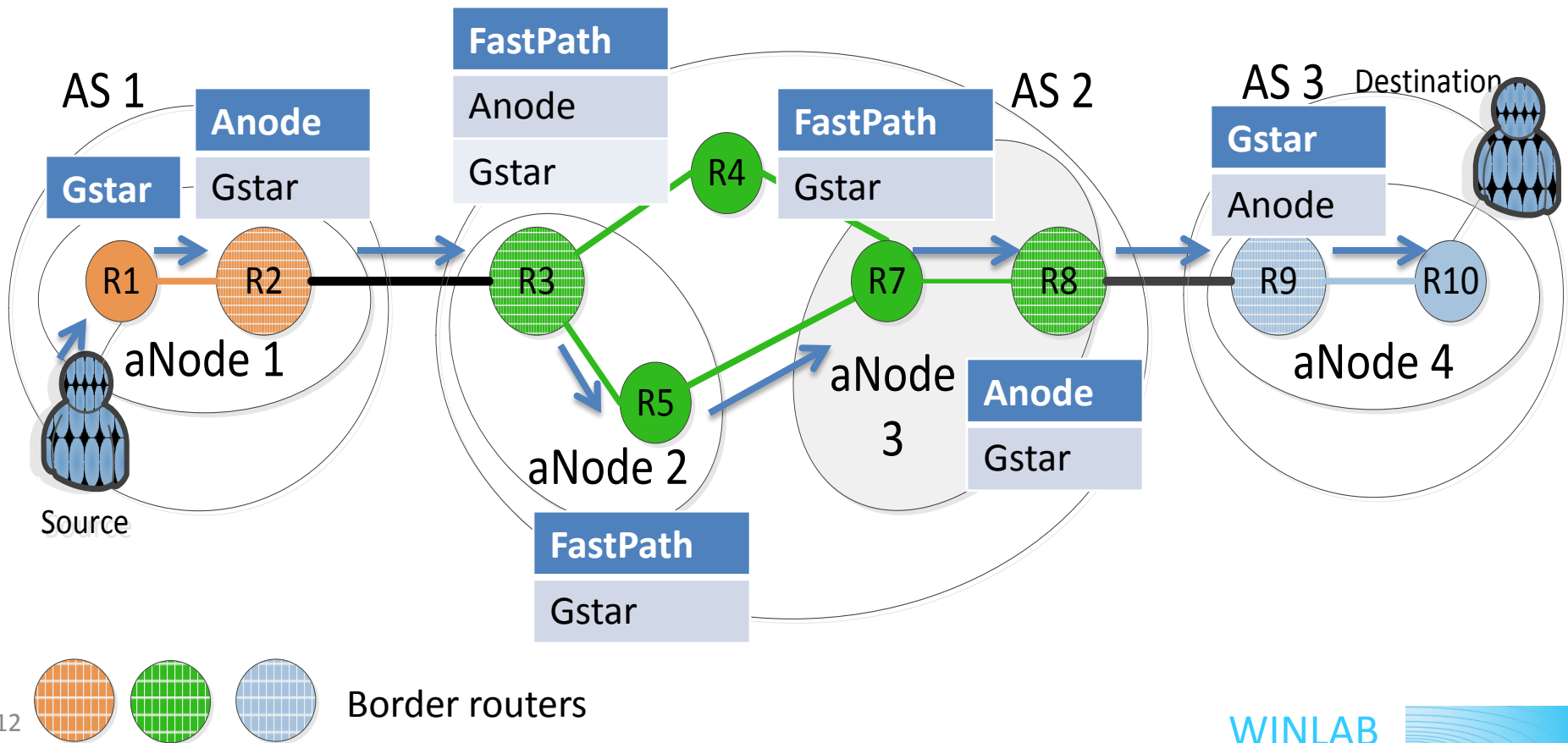


Border routers



# EIR Routing Scenarios

## 2. Source in AS 1 and destination in AS 3



# Implementation

- Working prototype implementation using Click modular routers on Orbit
  - Validated system performance for small to medium topologies
- Simulation implementation using NS-3
- Custom event driven simulator to simulate Internet scale topologies
  - Ongoing analysis of currently available topology datasets (Dimes, Caida) as well as future Internet topology generators (GeoTopo<sup>1</sup>)

<sup>1</sup> Y. Hu, F. Zhang, K.K. Ramakrishnan and D. Raychaudhuri, “GeoTopo: A PoP-level Topology Generator for Future Internet Study”, WINLAB Technical Report 2014



# Evaluation Plan

- Analysis of mobile wireless use-case scenarios
  - End users with realistic dynamic mobility
  - Cross-domain multi-homing
  - Wireless edge peering
- Scalability and control overhead evaluation for large scale Internet topologies



# Questions?

Thank you!

