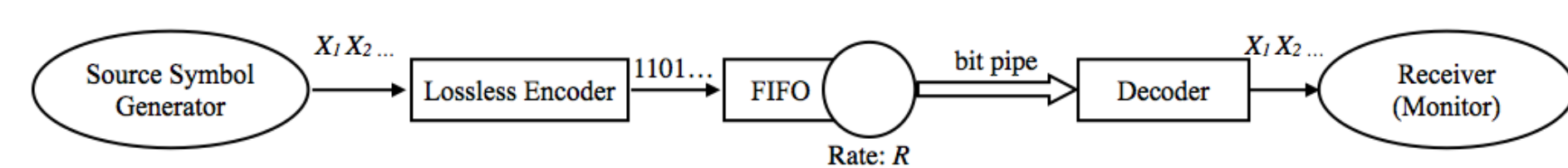


## Project Outlines

- We examine lossless data compression from an average delay perspective.
- Status updating system model:
  - A source generates time-stamped status update messages that are transmitted through a communication system to one or more receivers.
  - Updates must be as timely as possible.
- **Timeliness metric:** status update age.
- **Real-time lossless source coding** is a problem of status updating.
  - We evaluate the average status age in lossless sequential block coding systems.
  - We connect average age to the source coding error exponent with delay.
  - We propose the age-optimal block coding scheme.

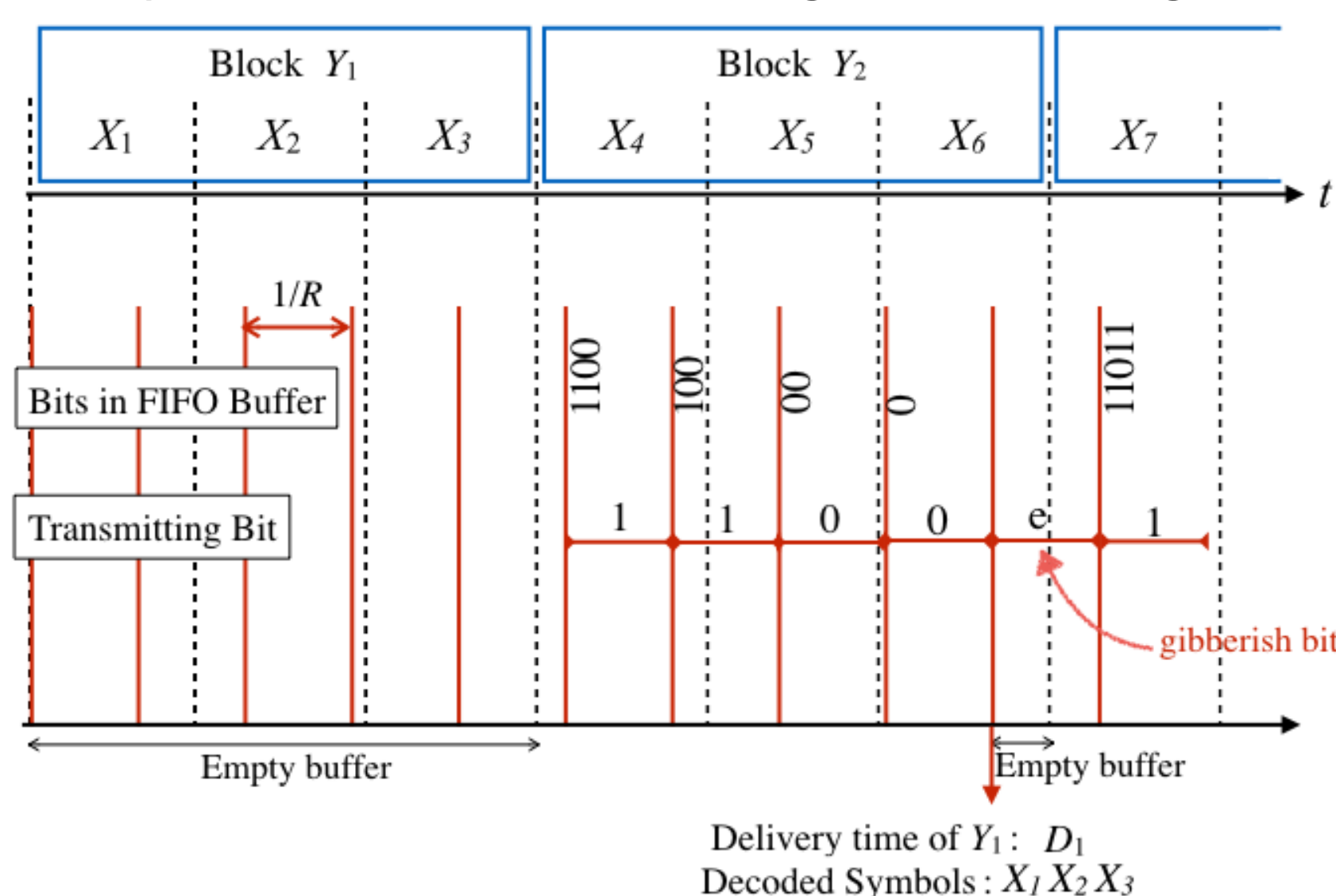
## System Model

- An encoder receives input symbols one per unit time from an i.i.d. source and submits binary codewords to a FIFO buffer that transmits bits at a fixed rate to a decoder.

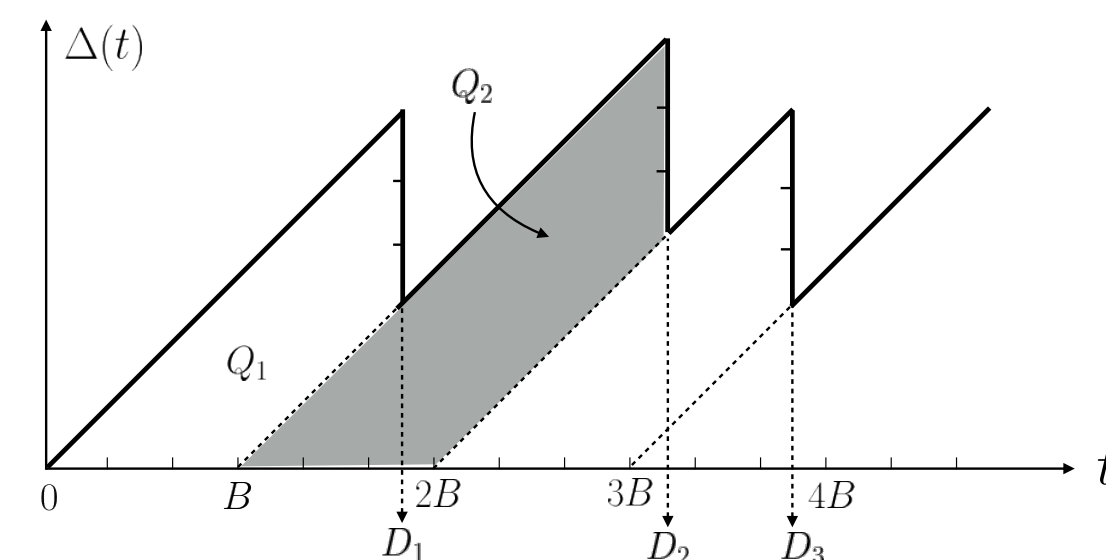


- Each input symbol at the encoder is viewed as a status update by the source.
- System performance is characterized by the **average status update age**:
  - The number of time units (symbols) the decoder output lags behind the encoder input.

Example of real-time block coding with blocklength  $B=3$



## Status Age Analysis



Example of status age for block coding schemes

- Average age for block coding:

- The average status update age is the normalized area under the sawtooth.

$$\Delta = \lim_{N \rightarrow \infty} \frac{1}{N} \sum_{k=1}^N (D_k - kB) + \frac{B}{2} = E[T] + \frac{B}{2}$$

average system time of symbol  $k$   $\leftarrow$  average waiting time + average service time

- D/G/1 queue model:

- service time is proportional to the length of the encoded sequence.
- the average waiting time can be solved numerically.
  - waiting times form a discrete Markov chain.
  - obtain the stationary distribution of M.C. from the discrete distribution of the service time.

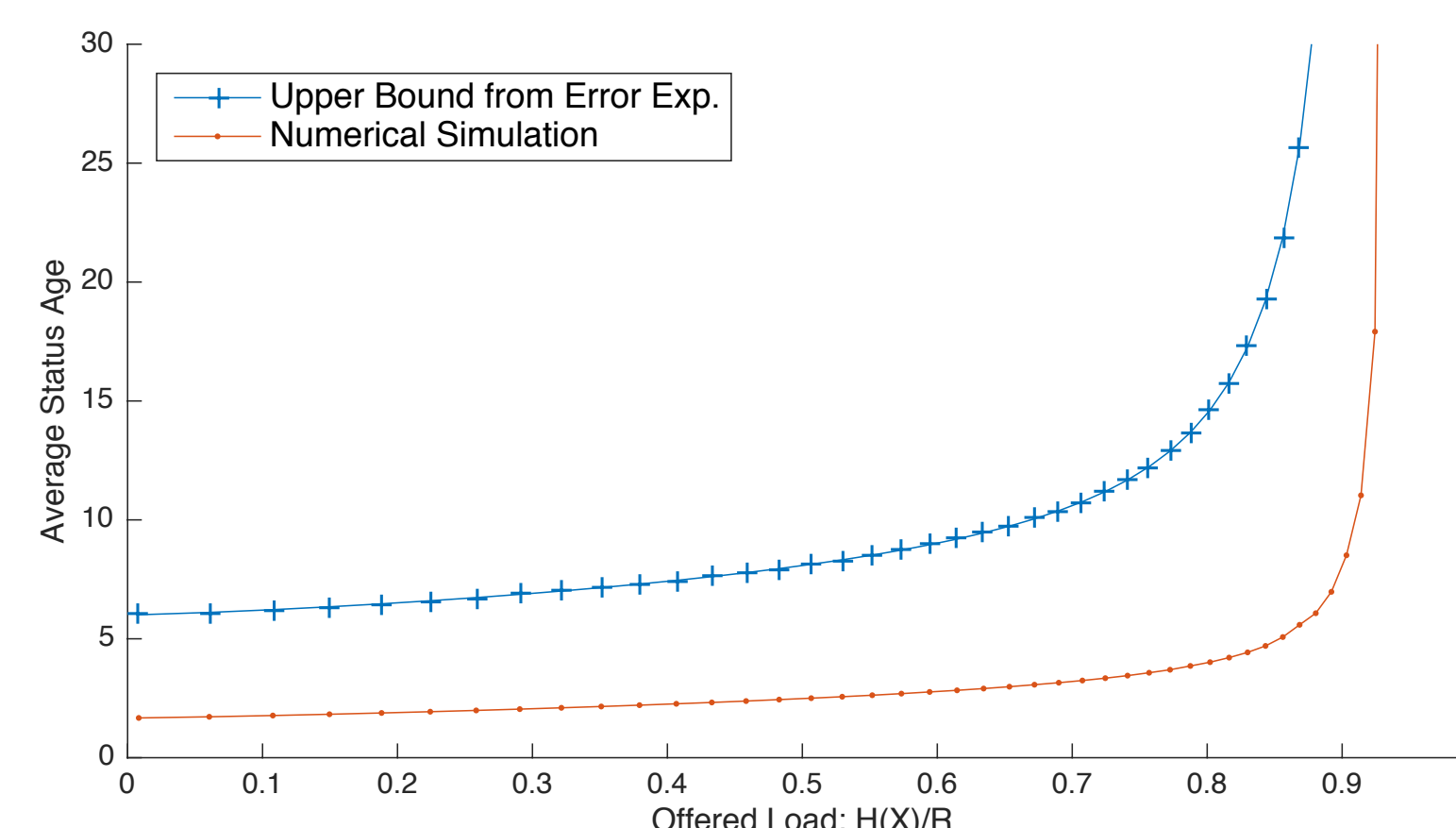
- Connection to source coding error exponent:

- a delay constrained error exponent is achievable if and only if

$$\Pr[\hat{x}_{n-\delta}(n) \neq x_{n-\delta}] \leq K 2^{(-\delta E_S(R) - \epsilon)}, \text{ for all } \delta \text{ and } n > \delta.$$

$$\Delta \leq \frac{K 2^{2E_S(R)}}{(2^{E_S(R)} - 1)^2} \triangleq f(K, E_S(R)).$$

- the error exponent describes the convergence rate of the error probability.
- this bound also scales with the constant term  $K$  in the error probability.



$$P_X(a) = p$$

$$P_X(b) = \frac{1-p}{2}$$

$$P_X(c) = \frac{1-p}{2}$$

Prefix-free block coding example with blocklength  $B=2$ :

- Encode  $aa$  to a single bit 0, and anything else to length-4 codewords.
- Fix channel rate  $R$  and vary the source entropy.
- The error exponent does not lead to an accurate description of the status age for small delay and small blocklength.

## Age-optimal Block Code

- Alternative upper bound on the average waiting time of G/G/1 queue:

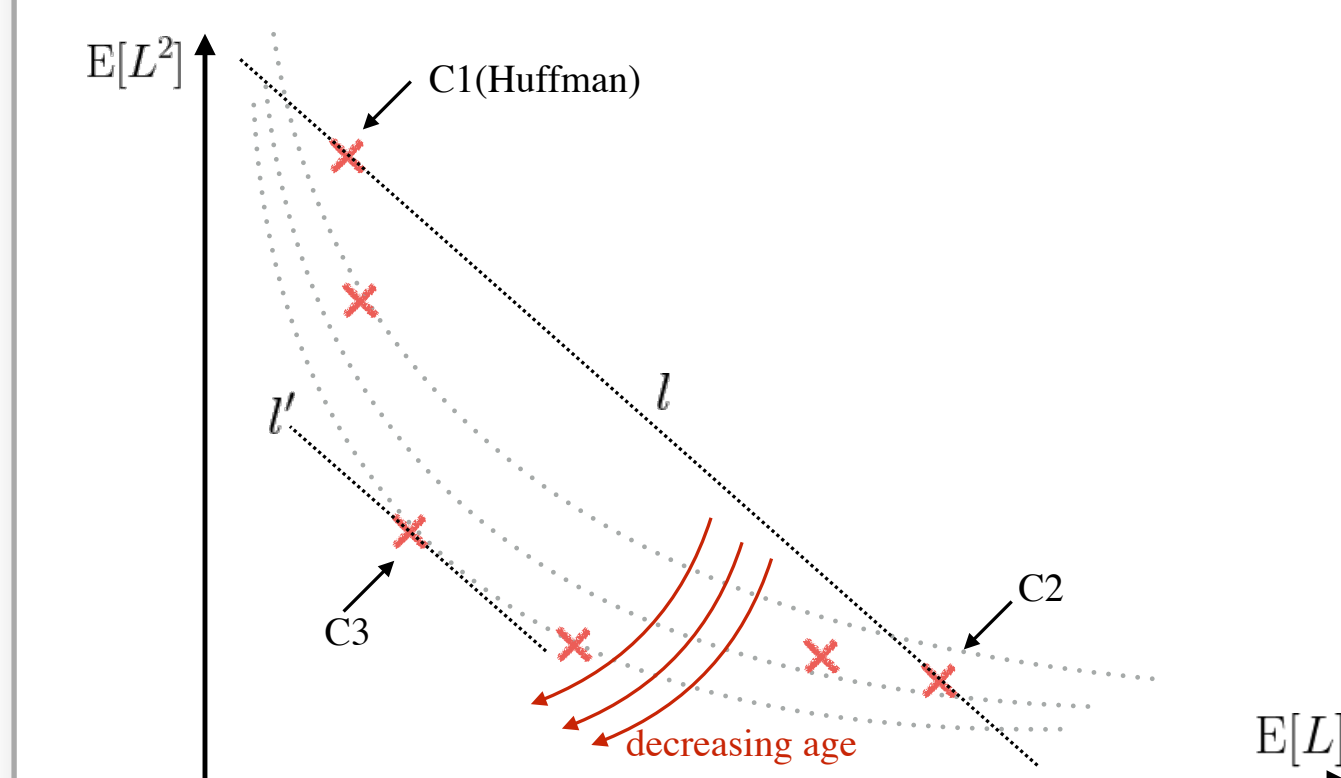
$$\Delta \leq \frac{\text{Var}[L]}{2R(BR - E[L])} + \frac{E[L]}{R} + \frac{B}{2}$$

$$= \frac{E[L^2] - E^2[L]}{2R(BR - E[L])} + \frac{E[L]}{R} + \frac{B}{2}$$

quasi-linear function of  $\leftarrow$  expected code length & second moment of code length

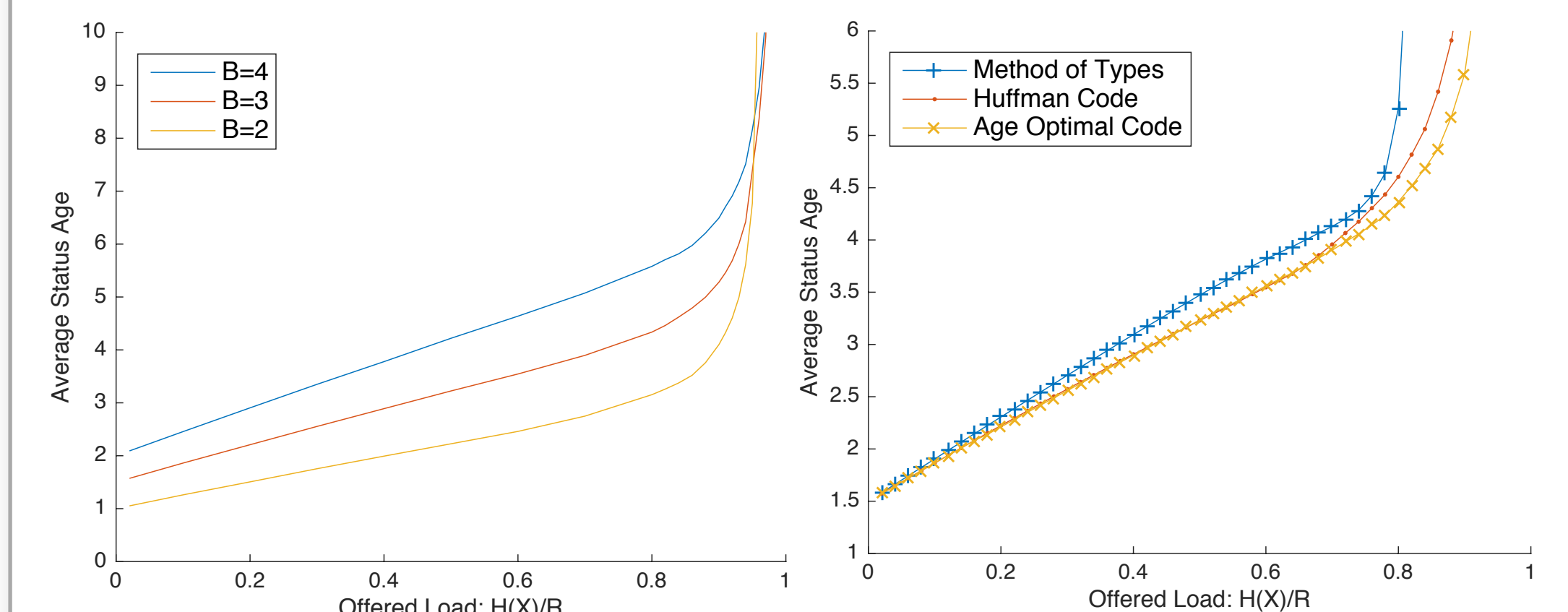
- Age minimized block code:

- search over all the possible codebooks lying on the lower left corner of convex hull iteratively.



The search process starts from two end points: Huffman code C1 and the minimum second moment code C2.

- Computes the analytical average age function iteratively and returns the optimal code tree that gives the minimum.



Comparison of different blocklengths:

- Encoding large blocks can result in high status age when the load is low.
- Small blocks encoding yields inefficient compression, which increases the offered load on the FIFO buffer and causes queueing delays.
- Error exponent is optimized by the scheme uses method of types.
- Huffman code is only optimal when the system is idle.
- With same blocklength, age-optimal code outperforms Huffman code significantly when the system is heavily loaded.

## Future Work

- Techniques to handle large blocklengths.
- Sources with memory.
- Age-optimal universal coding scheme.
- Coding with shared network resources.