

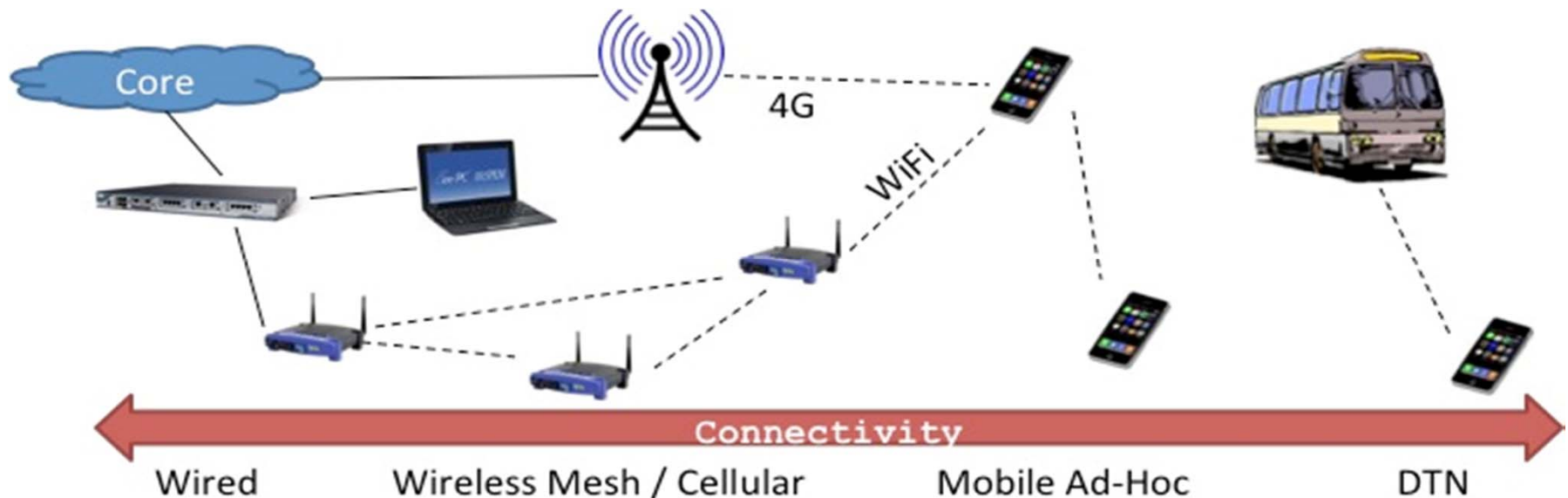
GSTAR: Storage Aware Routing Protocol for Efficient and Robust Services

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Routing in MobilityFirst: Objectives

- Efficient and robust support of mobility services in the core Internet
- Unified approach for handling all the challenges associated with mobile devices and associated applications
- Capable of achieving high performance across a wide range of wireless and wired networks



Routing in MobilityFirst: Approach

- Challenges associated with mobility are not addressed by current local-scale routing protocols.
- Some solutions have been proposed by:
 - Delay Tolerant Networking (DTN) Community –
 - Uses message replication and hop-by-hop transport
 - Not sufficient in highly connected environments
 - Ad-hoc/MANET Community –
 - Not sufficient in highly disconnected environments

Merging MANET and DTN

Generalized Storage Aware Routing (GSTAR)

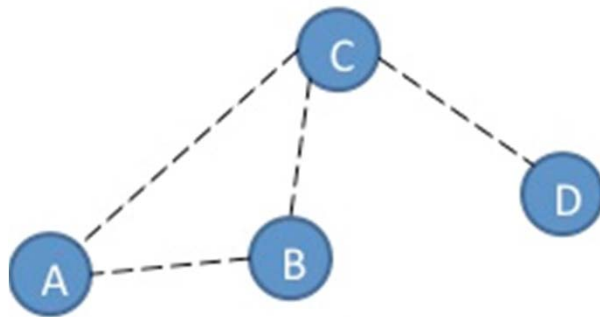
- Proactive link-state protocol with DTN capabilities for use in MobilityFirst networks
- Unifies techniques from MANET and DTN protocols
- Handles mobility related challenges at network layer using:
 - Exposed path quality information
 - Exposed connectivity patterns
 - Directly accessible in-network storage

GSTAR: An Overview

- Intra-partition Graph
 - Contains fine-grained, time sensitive information about the links
 - Uses Expected Transmission Time (ETT) as a measure of link quality
- Inter-partition Graph
 - Contains coarse-grained, time insensitive information about the connection probabilities
 - Based on Average Availability (AA) of nodes in the network
- Routing decisions are made on a set of data packets called *chunks*.

Intra-partition Graph: Control Messages

- Link Probe (LP)
 - Enables a node to know about the ETT of current one-hop neighbors
 - Used to compute short term expected transmission time (SETT) and long term expected transmission time (LETT)
- Flooded Link State Advertisement (F-LSA)
 - Contains SETT and LETT for all one-hop neighbors
 - Periodically flooded and re-transmitted by every node



D sees:

$IP_A \leftrightarrow IP_C$ (fair/fair)

$IP_A \leftrightarrow IP_B$ (good/fair)

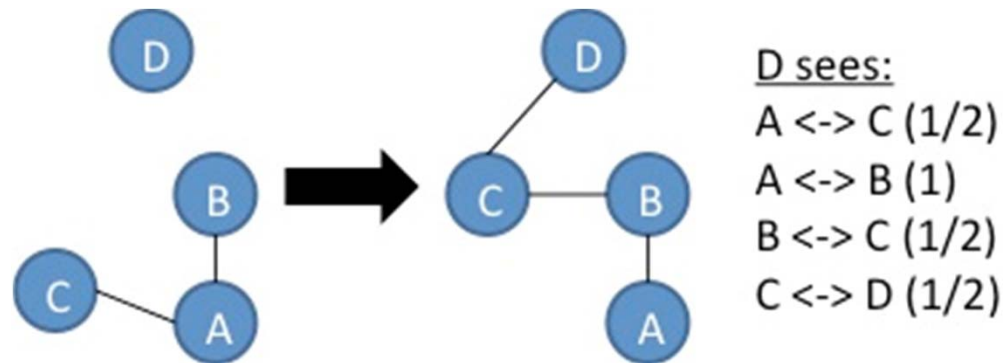
$IP_B \leftrightarrow IP_C$ (fair/bad)

$IP_B \leftrightarrow IP_C$ (good/good)

$IP_C \leftrightarrow IP_D$ (fair/fair)

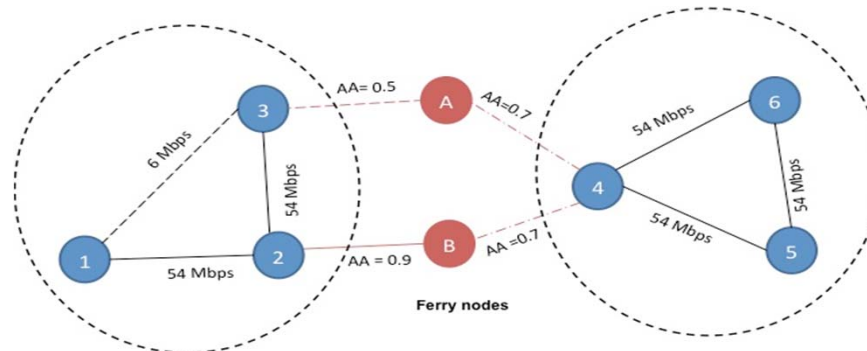
Inter-partition Graph: Control Messages

- Link Probe (LP)
 - Used to compute Average Availability (AA) as: $AA = \frac{on}{on + off}$
 - “on” time: active connection and “off” time: disconnection
- Disseminated Link State Advertisement (D-LSA)
 - Contains AA for all nodes in the complete network
 - Epidemically disseminated and carried indefinitely by every node



Intra-partition Forwarding Table

- Computed using any shortest path algorithm like Dijkstra's with SETT as link weights
- Contains only *end-to-end routes* with the corresponding SETT and LETT



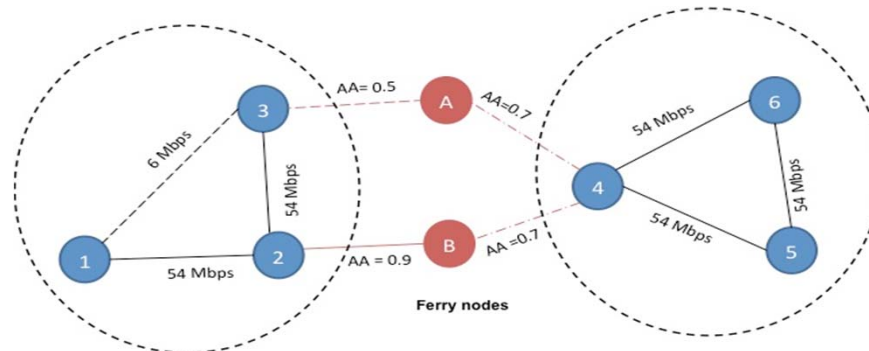
Intra-partition table at Node 1

Dest	Next Hop	ST Path	LT Path	Hops
2	2	13332	13332	1
3	2	66666	66666	2

ST Path – SETT Sum and LT Path – LETT Sum

Inter-partition Forwarding Table

- Computed using any shortest path algorithm like Dijkstra's with link weights as: $(1-AA+0.01)$
- Contains *highly probable routes* to all nodes in the network



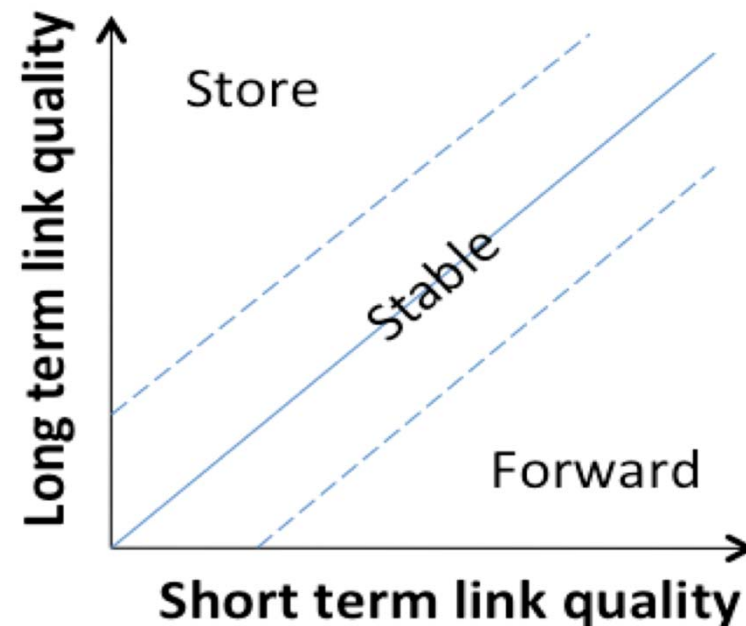
Intra-partition table at Node 1

Dest	Next Hop	AA	Dest	Next Hop	AA
2	2	0.01	4	2	0.43
3	3	0.01	5	2	0.44
A	3	0.52	6	2	0.44
B	2	0.12			

Transmission of Data

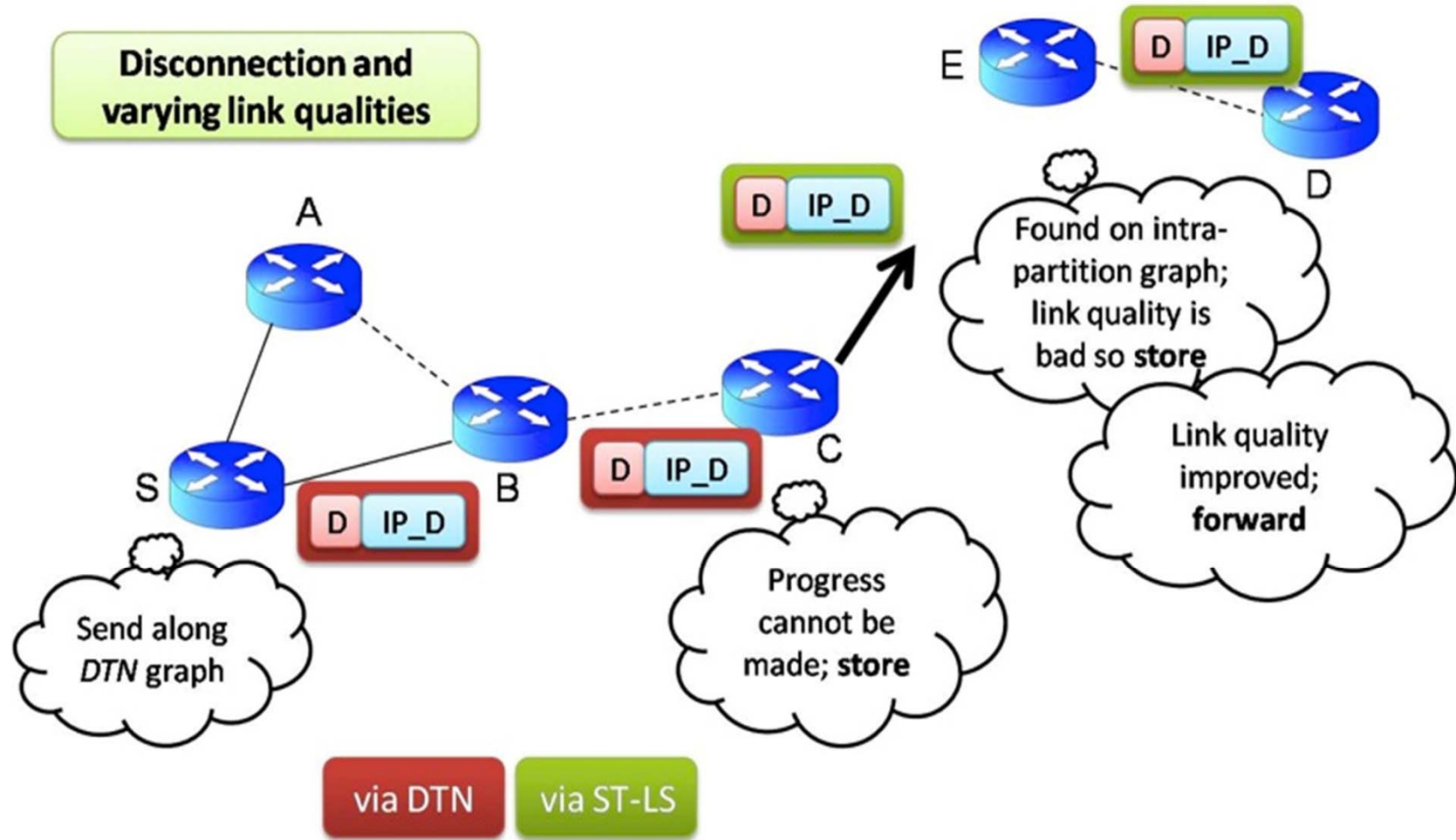
- A node first checks its intra-partition table for an end-to-end route to the destination.

- if ($SETT > 1.1 * LETT$)
 store the chunk
else
 forward



- If no route exists in the intra-partition table, the node switches to DTN mode and checks the inter-partition graph.

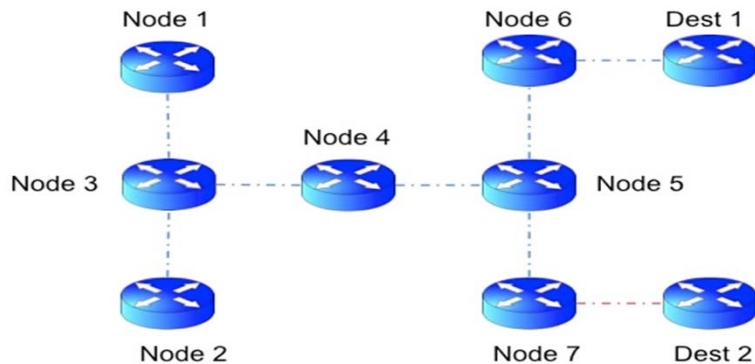
Working of GSTAR



Simulation Model

- NS-3 (Network Simulator 3) based simulation model is developed for evaluation of GSTAR.
- The simulation model consists of:
 - Nodes with storage
 - Hop-by-hop transport
 - Time varying wireless channel
 - Mobile users with possible disconnection

GSTAR vs. Link State in Wireless Network

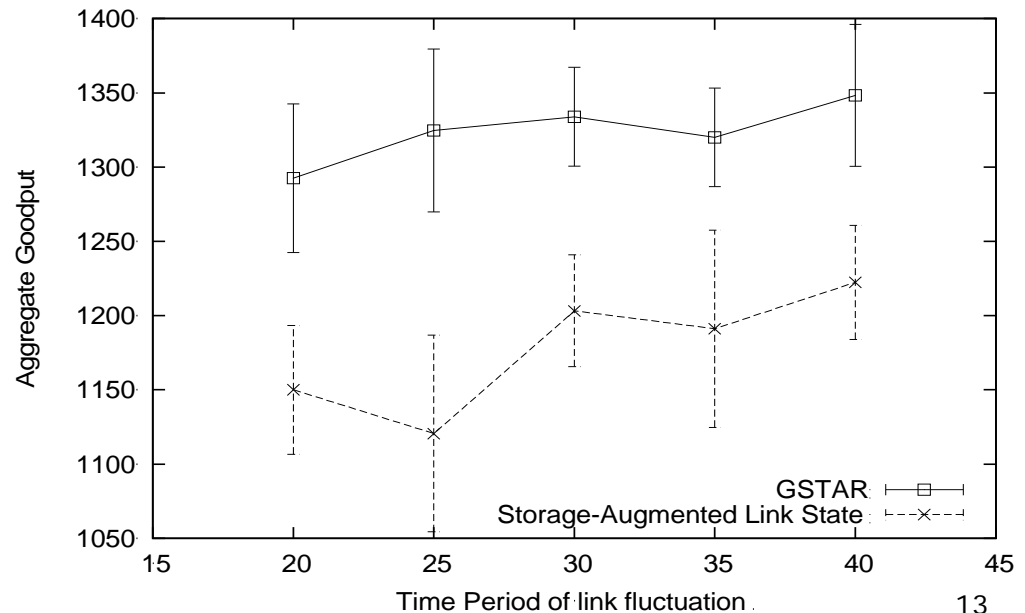


Simulation Parameters

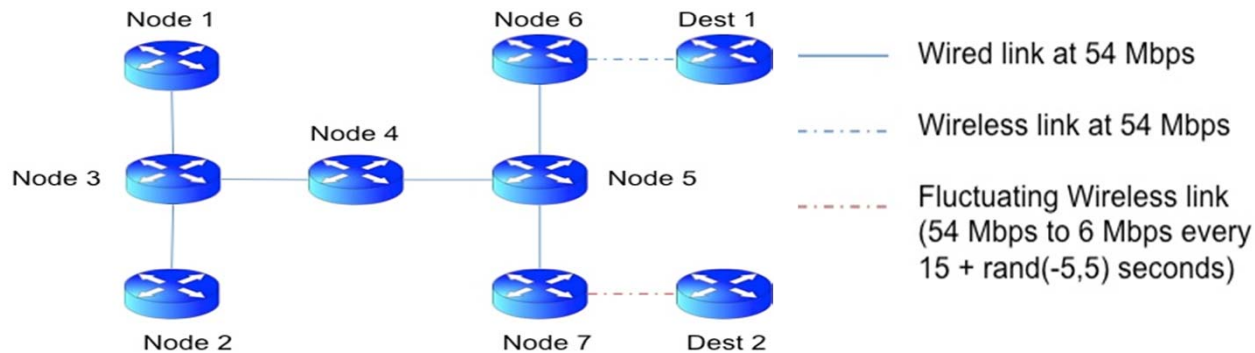
- Flows: Node 1 – Dest 1
Node 2 – Dest 2
- Chunk Size: 25 packets
- Simulation Time: 90 sec
- Each data point is average of 10 runs

- LETT: average of past 10 ETTs
- Store-forward decision threshold: 1.1

GSTAR alleviates the effect of congestion in Flow 2 from Flow 1 resulting in better network utilization.



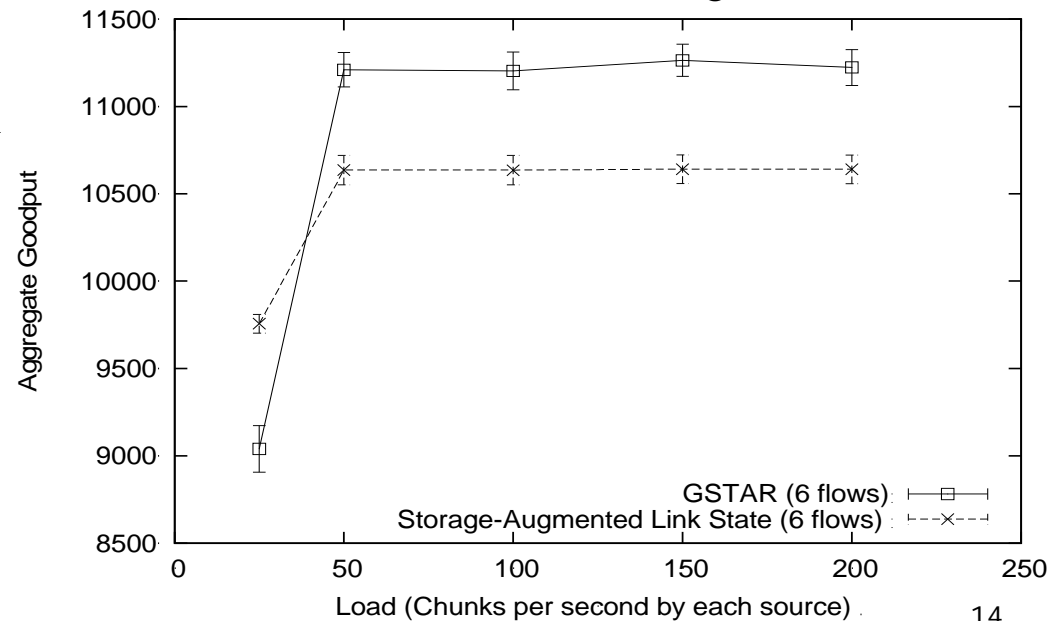
GSTAR vs. Link State in Hybrid Network



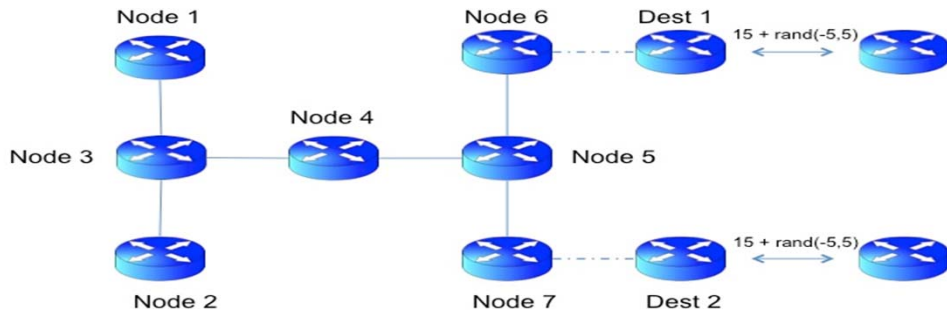
Simulation Parameters

- Flows: 3 flows to Dest 1
3 flows to Dest 2
- Chunk Size: 10 packets
- Simulation Time: 90 sec
- Each data point is average of 10 runs

- LETT: average of past 10 ETTs
- Store-forward decision threshold: 1.1
- GSTAR provides a gain in aggregate goodput for medium to high offered load.**
- Cross-over point is the load at which network is fully utilized.**



GSTAR w/ DTN vs. GSTAR w/o DTN

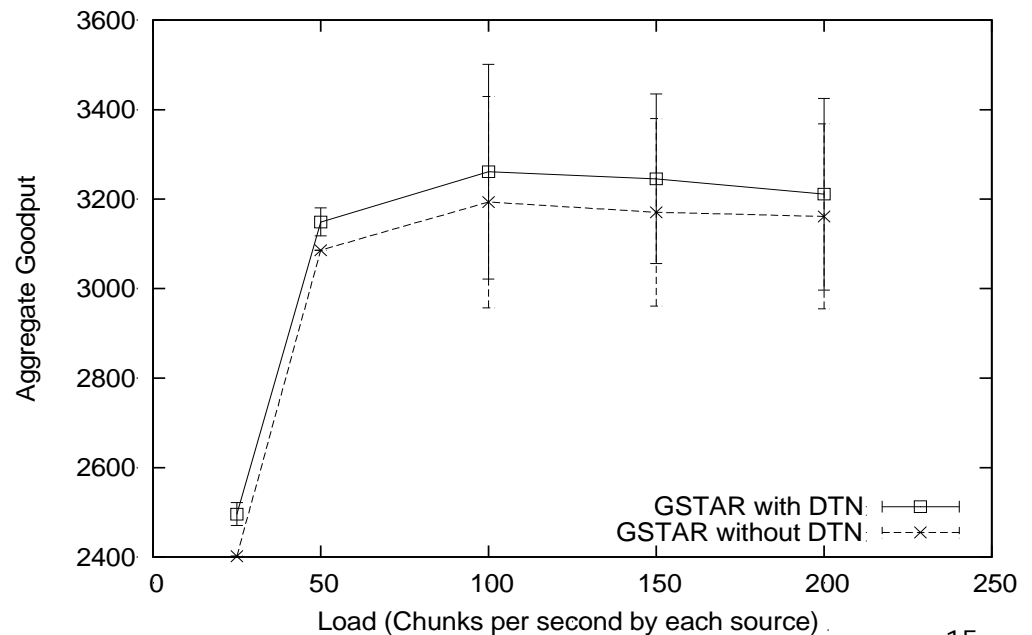


— Wired link at 54 Mbps
 - - - Wireless link at 54 Mbps

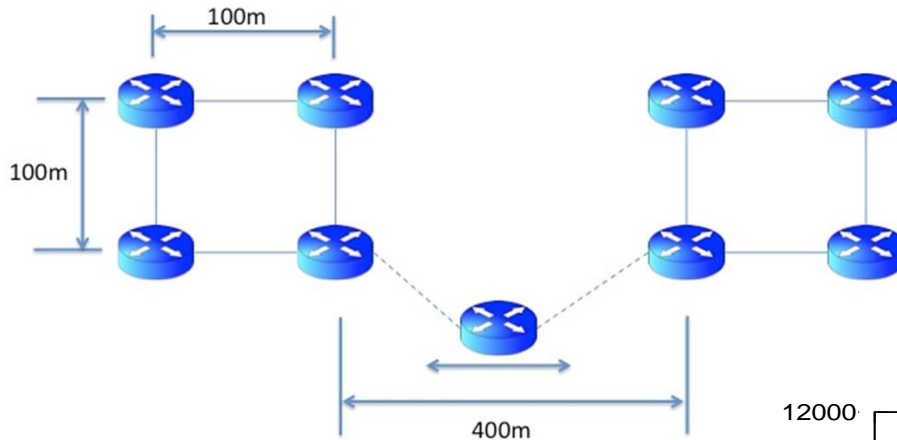
- **Proactive pushing enables destinations to start receiving data as soon as it reconnects.**
- **W/o DTN, the destinations have to wait for there F-LSAs to be received by the sources.**

Simulation Parameters

- Flows: Node 1 – Dest 1
Node 2 – Dest 2
- Chunk Size: 25 packets
- Simulation Time: 90 seconds
- Each data point is average of 10 runs



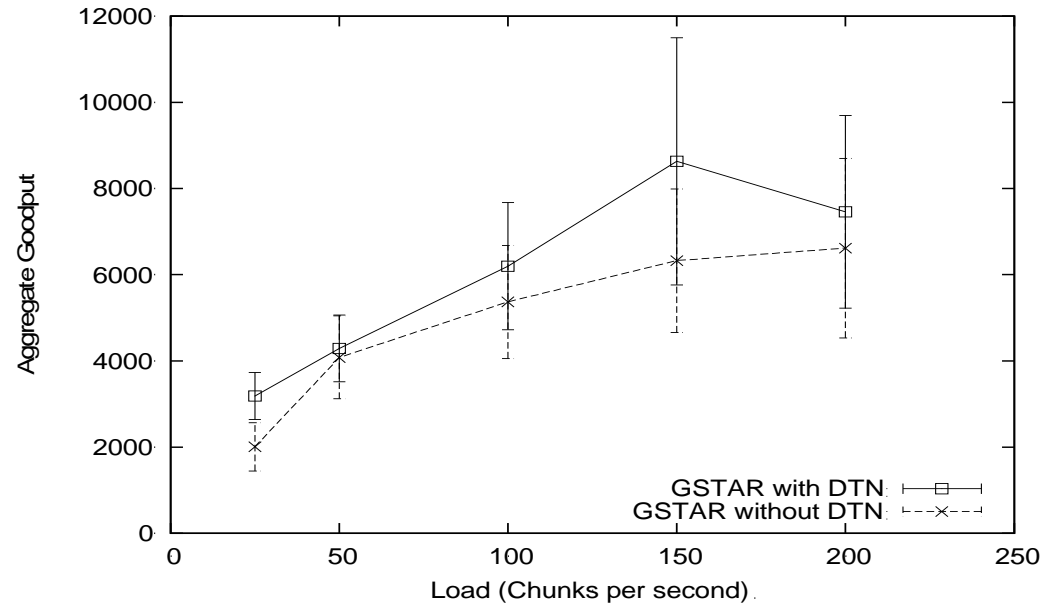
GSTAR vs. Link State with Network Partitions



Proactive pushing enables data to be received across network partitions.

Simulation Parameters

- Flows: 3
- Chunk Size: 25 packets
- Simulation Time: 120 seconds
- Each data point is average of 10 runs



Computation of LETT

1) Exponentially weighted moving average (EWMA)

- $LETT = \alpha \cdot SETT + (1 - \alpha) \cdot LETT$
- α is the weighting factor explored via simulation
- Works well for periodic links as past information is relevant

2) Simple moving average

- Giving equal weights to past ETTs
- Using different amounts of past ETTs is explored via simulation
- Works well if the link fluctuation period is less than the amount of past history used

Adaptive Store-Forward Decision Threshold

- Static threshold of 1.1 works well
 - With simple on-off model
- Adaptive or dynamic threshold works well
 - For networks where link fluctuation model is unknown
 - Approaches-
 - **Simple Moving Average Filtering:** Average of past *ten* SETT/LETT ratio
 - **Median Filtering:** Median of past *ten* SETT/LETT ratio
 - **Moving Average + Median Filtering:** Perform averaging of past *five* SETT/LETT ratio and then median filtering on *five* such averaged ratio values

Future Work

- Inter-partition Graph
 - Comparing current single copy DTN routing to multiple copy DTN routing mechanisms
 - Comparing GSTAR with existing DTN routing protocols
- Storage Aware Routing Metric
 - The path selection metric should be modified to include SETT, LETT and storage available at each router.
- Effects of finite storage at each router
- Extending GSTAR to support multicast and anycast



Thank you !