

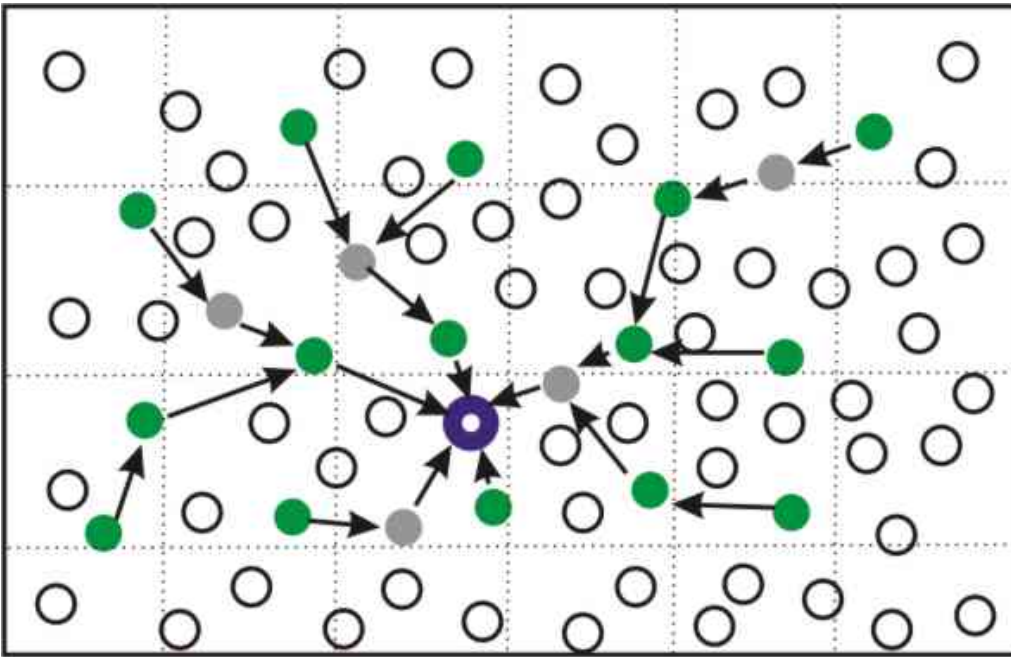
---

# PROSE: Providing Robustness in Systems of Embedded Sensors

*Yanyong Zhang, Shengchao Yu, Jaewon  
Kang, Badri Nath, Wade Trappe*

WINLAB, Rutgers University

# Sensor Network Model

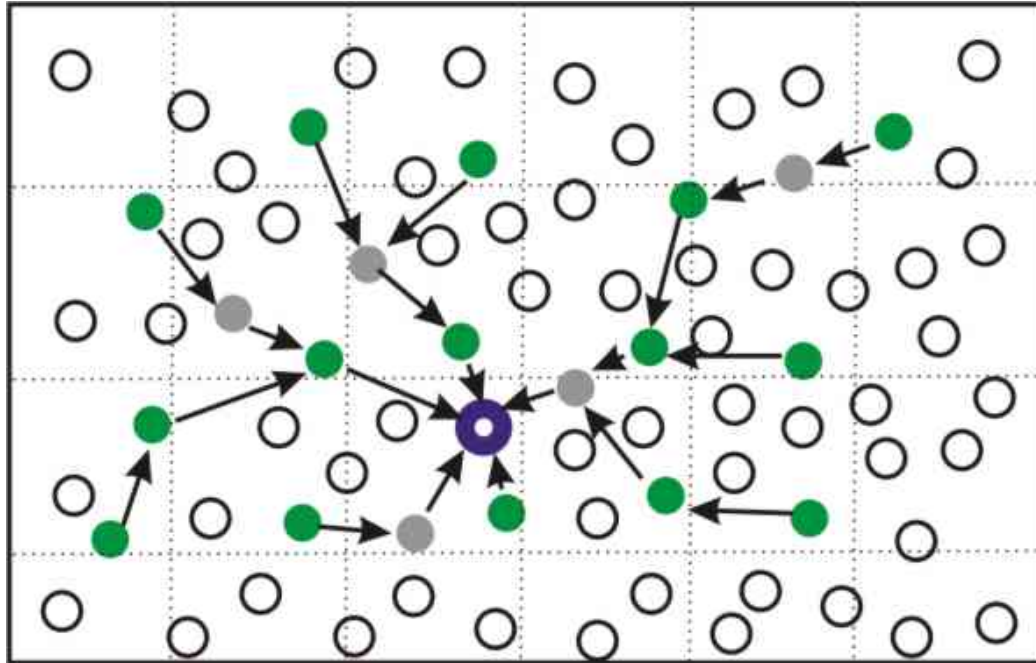


## □ Sensor net connectivity model

- Multi-hop networking
- More active nodes may be needed

- A sensor net deployment has:
  - Many inexpensive sensor nodes
    - capable of sensing, computing, and communication
    - Resource constraints (energy)
    - Deployment redundancy
    - Stationary nodes
  - One or several data sinks
    - Connected with apps
- Sensor net coverage model:
  - Lifetime vs. coverage
    - a minimum set of active nodes
  - Grid-based coverage model
    - Network partitioned into grids
    - A node is able to monitor at least one grid point
    - Every grid point must be monitored

# Challenges in Sensor Networks



□ Many challenges in building sensor systems/applications (at WINLAB)

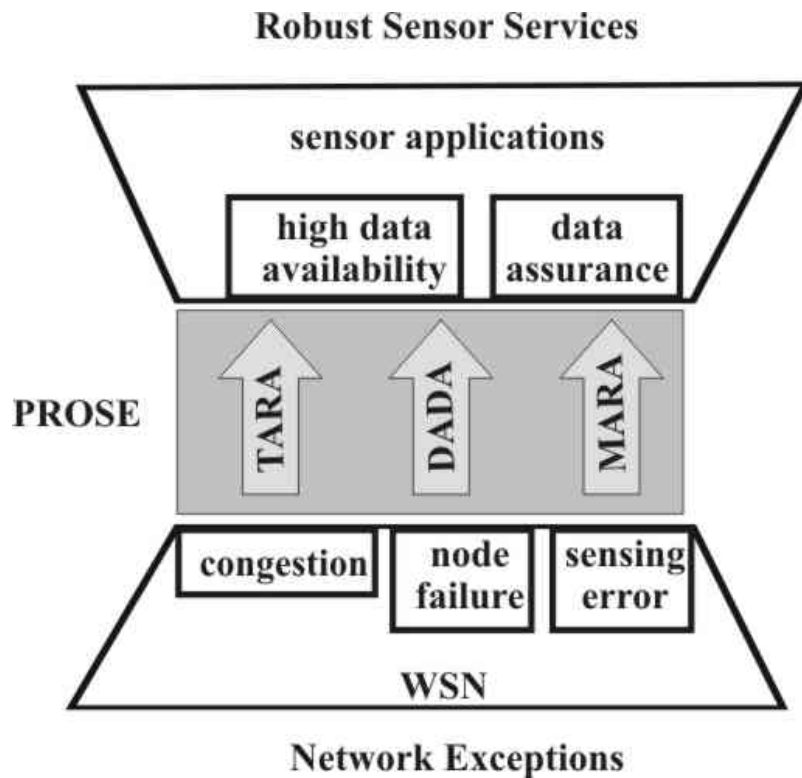
- Sensor deployment
- Lifetime maximization
- Programming model
- Routing and MAC protocols
- Localization algorithms
- Security and privacy

□ This project focuses on **robust** sensor services

- Unattended operations
- Hostile environment
  - Fire
  - landfill
- Many unexpected events may happen
  - Node failure
  - Congestion
  - Erroneous data

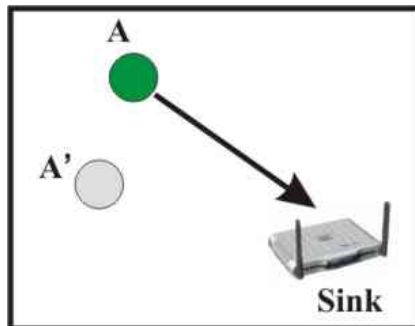
How to extract a steady stream of valid data from all the interesting spots in the presence of unexpected interrupts?

# Prose Overview

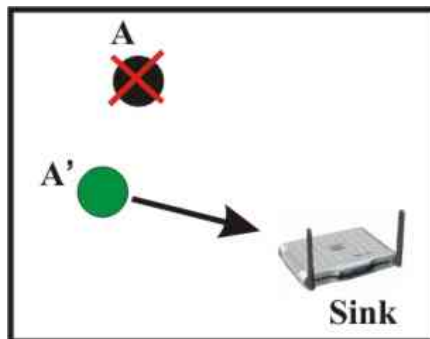
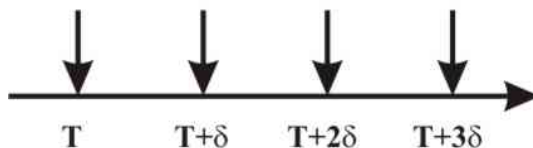


- PROSE has three components:
  - **DADA**: a 2-Dimensional Adaptive Node Scheduling Framework
    - Provides data availability against random node failures
    - Repairs network coverage and connectivity by cleverly waking up redundant nodes upon node failures
  - **TARA**: a Topology-Aware Resource Adaptation Framework
    - Provides data availability against congestion
    - Brings more sensor nodes "online" to accommodate higher traffic rate to eliminate congestion
  - **MARA**: a Measurement Assurance and Robust Aggregation Framework
    - Provides data assurance
    - Classifies and cleanses sensed data
- Each component has two units:
  - Exception detection
  - Exception handling

# DADA: Overview



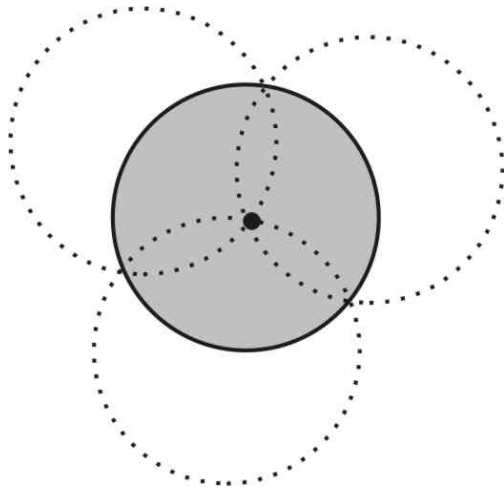
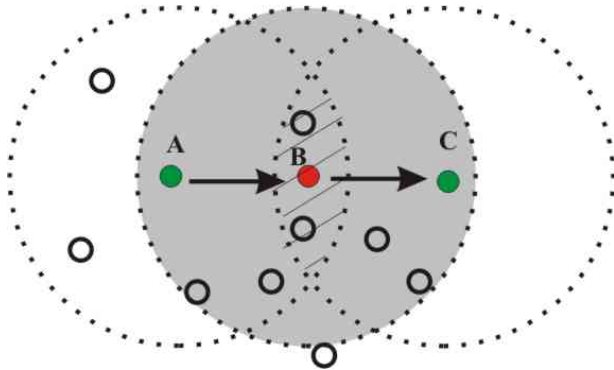
wake up schedule for A'



- DADA balances the tradeoff between network lifetime and network quality
- DADA achieves two goals:
  - Bounded recovery time upon node failures (recovery time  $\leq \delta$ )
  - Minimize energy consumption
- Basic idea for DADA
  - A minimum set of active nodes
  - The redundant node wakes up every  $\delta$  time
  - When it wakes up, it finds out whether it needs to become active
  - If the active node dies at time  $t$ , its redundant node will wake up at latest  $t+\delta$  to, and will become active
- Complications
  - How to decide an active node's redundant nodes?
  - How to schedule multiple redundant nodes?

# DADA: Gangs

---



- ❑ A node can not always be replaced by one redundant node
  - On average, 3-5 nodes are needed to replace a node's sensing area
  - One node may not be enough to repair the communication hole even if it is in the radio range
- ❑ The concept of "gangs"
  - A **gang** consists of a group of nodes that can completely replace an active node
  - A **minimum gang** is a gang itself, but none of its subsets is a gang
  - All the nodes belong to a minimum gang should wake up together
  - The minimum gangs are analogous to "sentries" in real life.

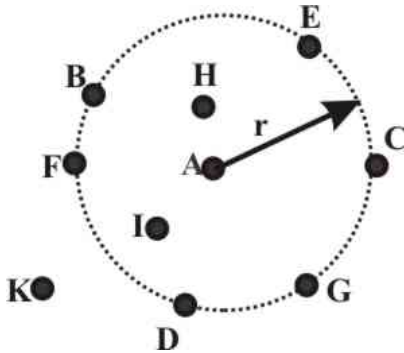
# DADA: P-Sentry Algorithm

---

- ❑ Persistent Sentry (P-Sentry) provides interrupt-less network operations
  - One minimum gang stays awake all the time (the sentry)
  - The sentries can take over when the active node fails
  - The other redundant nodes can sleep for a much longer time
- ❑ Issues with P-Sentry algorithm
  - How do the sentries detect the failure of the active node?
    - Passive listening
    - Active probing
  - Which nodes should be chosen as sentries?
    - Energy
    - Functionality
    - Redundancy
  - How long should the non-sentry redundant nodes sleep?
    - Estimate the remaining lifetime
    - Considering the random failure rate
  - How to synchronize multiple schedules a redundant node may have?
    - Active node maintains states



# DADA: R-Sentry Algorithm



- Rotary sentries (R-Sentry) limits service loss

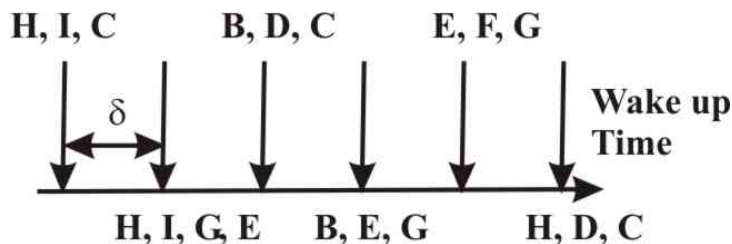
- All the minimum gangs take turns to wake up, with  $\delta$  between two subsequent wakeups.

- An example scenario

- A's minimum gangs  $\{H, I, C\}$ ,  $\{H, I, G, E\}$ ,  $\{B, D, C\}$ ,  $\{B, E, G\}$ ,  $\{E, F, G\}$ , and  $\{H, D, C\}$

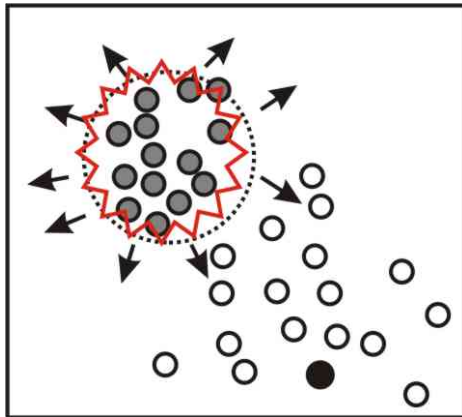
- Issues to be considered

- How to detect failures?
  - Probing for sensing ranges
  - Listening for HELP messages
- How to synchronize multiple schedules a redundant node may have?
- How to dynamically adapt the schedule when the redundant node fails?

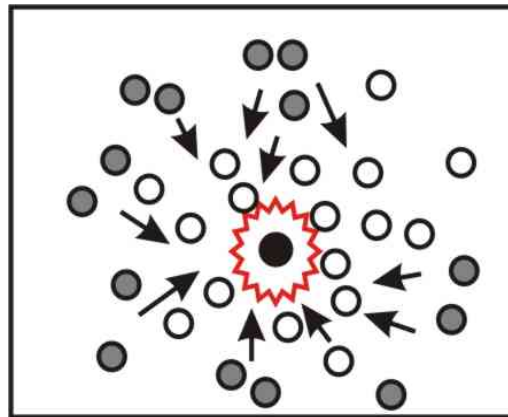




# TARA: Overview I

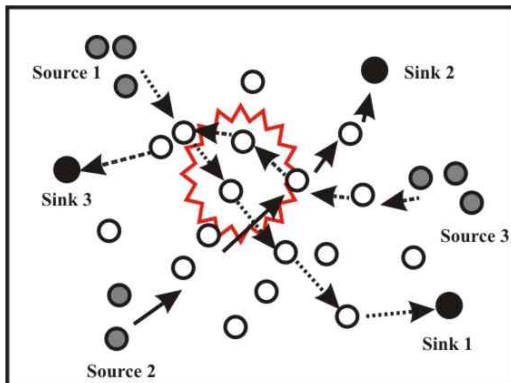


Source hot spot

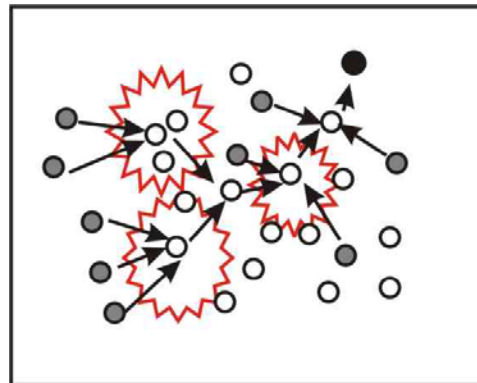


Sink hot spot

- A monitoring sensor network alternates between dormant periods and crisis periods
  - During dormant periods, minimum resources are kept online (DADA)
  - A hot spot will form during crisis periods



intersection hot spot



- Three types of hot spots:
  - Source hot spot
  - Sink hot spot
  - intersection hot spot

# TARA: Overview II

---

- ❑ Traffic control vs. resource control
  - Traffic control schemes throttle source traffic rates to eliminate hot spots
  - Resource control is preferred in sensor networks
    - Data packets during crisis states can NOT be dropped
      - Fidelity requirement
    - Sensor nets are deployed for peak load
    - Bringing resources online is realistic
      - e.g. sensor nodes
- ❑ Resource control schemes
  - Forming alternative routing topology
  - Multiple-path routing
- ❑ Topology-aware resource control
  - e.g. dose the new topology provide sufficient capacity?
- ❑ TARA has two main components:
  - Capacity estimation tool
  - Resource adaptation algorithm

# TARA: Capacity Analysis

---



Capacity fraction =  $1/2$

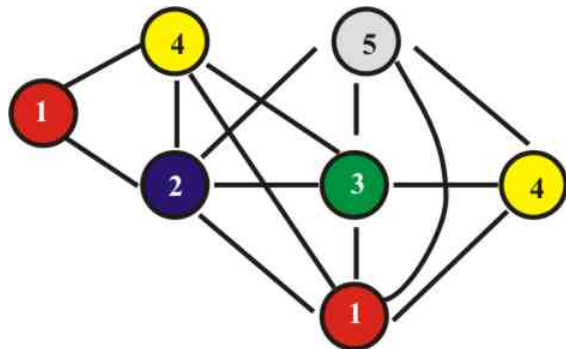
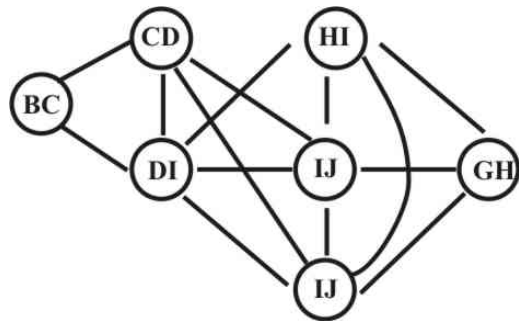
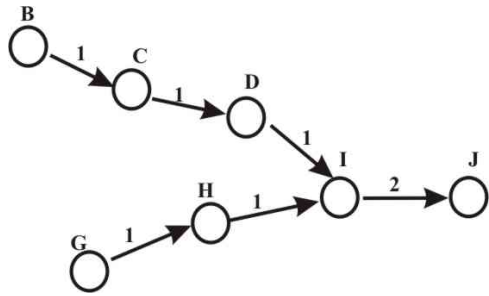


Capacity fraction =  $1/3$

Assumption: only adjacent nodes can hear/interfere each other

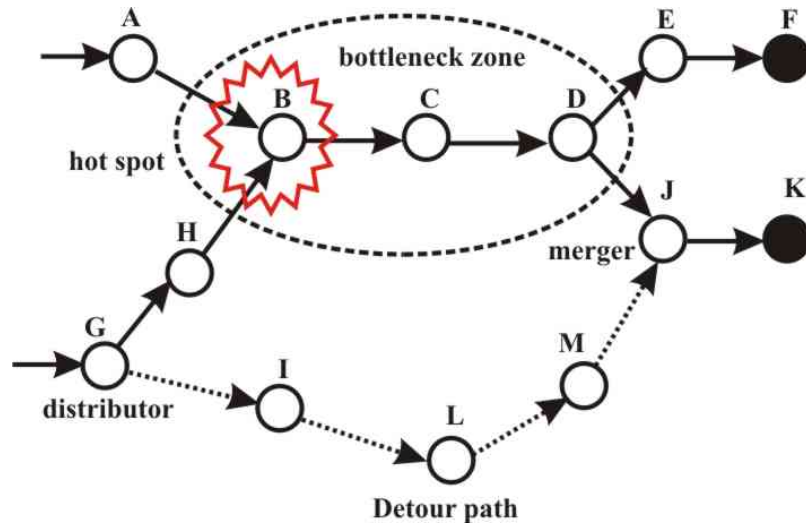
- Goal: to estimate the maximum end-to-end throughput of a given topology
  - The maximum rate at which the source can send packets towards the sink
- Idea: mapping this problem to a graph coloring problem
  - Due to link interference, the end-to-end throughput is a fraction of the 1-hop throughput (capacity fraction)
  - Suppose under an optimal schedule, the sink receives a packet every  $n$  time frames, then the capacity fraction is  $1/n$
  - The throughput is (1-hop throughput \* capacity fraction)

# TARA: Capacity Analysis



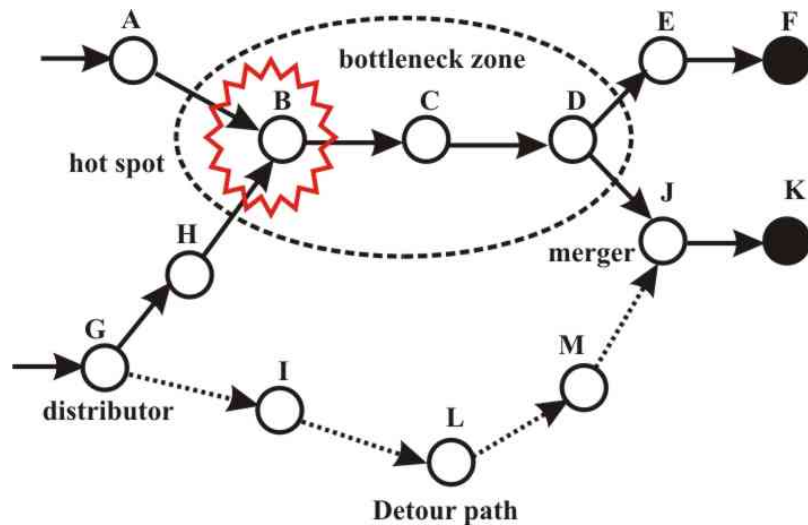
- ▣ Idea: mapping this problem to a graph coloring problem
  - Suppose under an optimal schedule, the sink receives a packet every  $n$  time frames, then the capacity fraction is  $1/n$
  - To estimate  $n$ 
    - ▣ Spatial interference graph
      - Vertexes are wireless links
      - Edges mean that two links are within each other's interference ranges
    - ▣  $N$  is equal to the number of colors assigned to the spatial interference graph
- ▣ Graph coloring problem
  - Theorems can provide an upper bound
  - Heuristic approaches provide tighter estimates

# TARA: Resource Adaptation Scheme



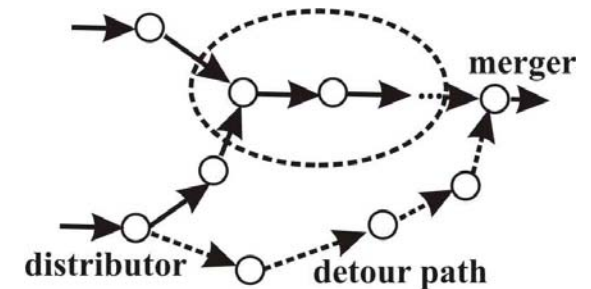
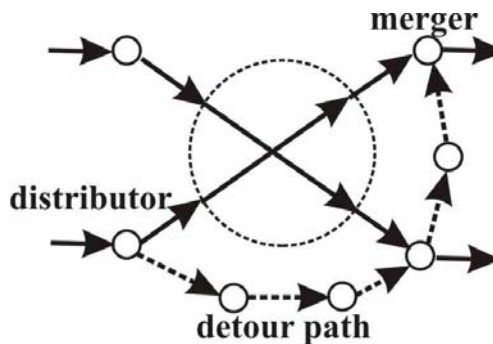
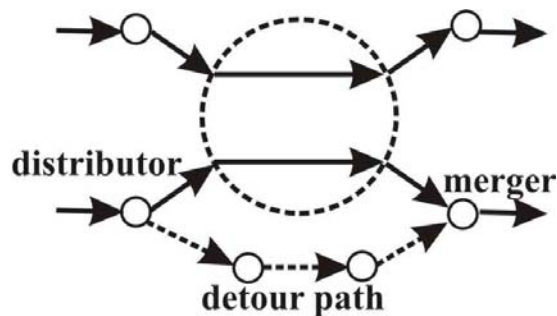
- TARA has the following steps:
  - Congestion detection
    - e.g. channel loading, packet drop ratio, queue occupancy, etc.
  - Traffic distributor
    - Each node maintains the incoming traffic rate from each neighbor

# TARA: Resource Adaptation Scheme



□ TARA has the following steps:

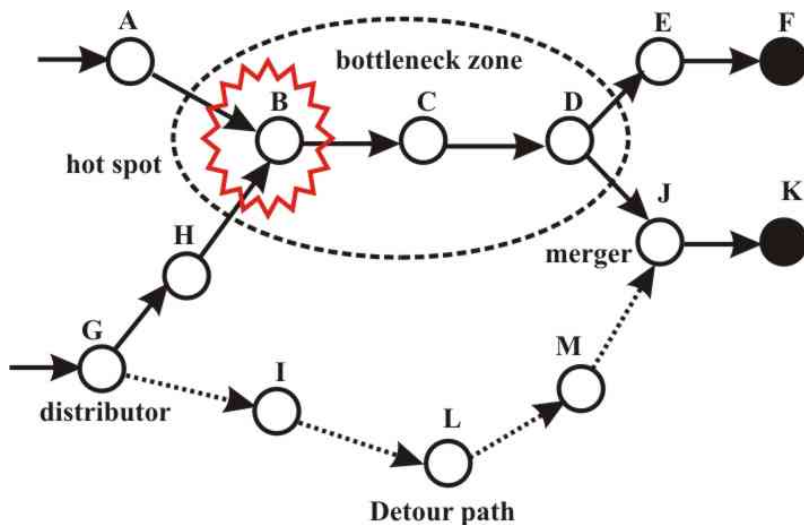
- Congestion detection
- Traffic distributor
- Traffic merger should
  - Have a low congestion level
  - Reside on the routing path to the intended sink
  - Result in sufficient capacity
    - Based on the observations gained using the capacity analytical tool



# TARA: Resource Adaptation Scheme

## □ TARA has the following steps:

- Congestion detection
- Traffic distributor
- Traffic merger
- Build a detour path
  - The merger initiates the process with flooding a REQ message including the TTL
  - A node that receives the REQ message decrements the TTL, and forwards the message with the following optimizations:
    - A node drops REQ if its congestion level is high
    - A node drops REQ if it is on the original routing path
    - A node only forwards the REQ with a higher TTL
  - Multiple REQ messages will reach the distributor, and it chooses the path with highest TTL

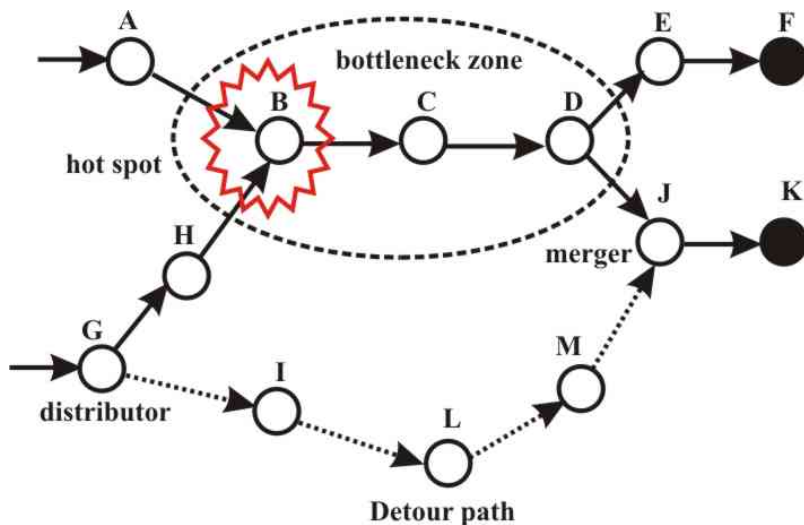




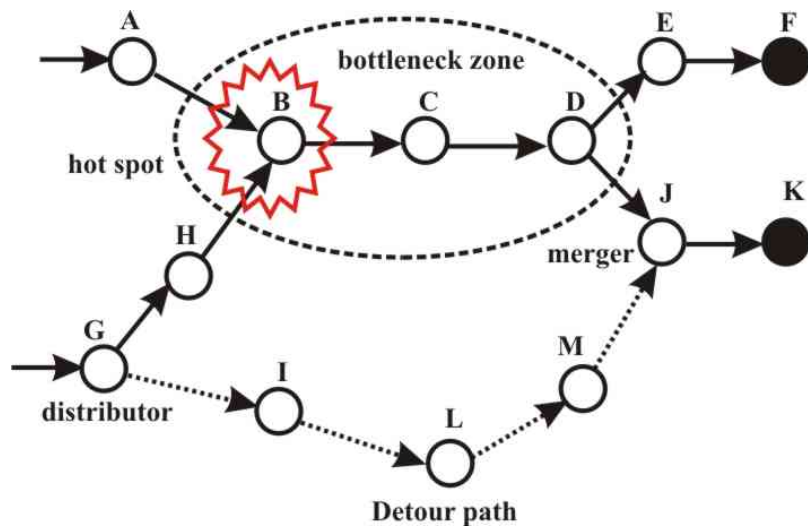
# TARA: Resource Adaptation Scheme

## □ TARA has the following steps:

- Congestion detection
- Traffic distributor
- Traffic merger
- Build a detour path
  - The merger initiates the process by flooding a REQ message including a TTL field
  - A node that receives the REQ message decrements the TTL, and forwards the message with the following optimizations:
    - A node drops REQ if its congestion level is high
    - A node drops REQ if it is on the original routing path
    - A node only forwards the REQ with a higher TTL
  - Multiple REQ messages will reach the distributor, and it chooses the path with highest TTL



# TARA: Resource Adaptation Scheme



- TARA has the following steps:
  - Congestion detection
  - Traffic distributor
  - Traffic merger
  - Build a detour path
  - The distributor distributes the traffic between two paths
    - Traffic allocation is topology-dependent

# TARA: A few issues ...

---

- ❑ TARA also reduces the online resource when traffic decreases
- ❑ A three-tiered resource controlling scheme
  - For very short-term congestion
    - Larger buffer size
    - Prioritizing packets
    - Storing data locally
  - For short-term congestion
    - TARA
  - For longer-term congestion
    - Traffic control

# MARA

---

- ❑ MARA provides data assurance against sensing errors
- ❑ Data classification mechanisms
  - Constraint-based Consistency Checks
    - Predefined constraints such as “temperature between 0 and 100”
  - Redundancy Consistency Checks
    - Multiple sensors monitoring the same variable
  - Multi-modal Consistency Checks
    - Multiple physical properties may exhibit correlation
- ❑ Data cleansing mechanisms
  - Robust statistical tool or robust aggregation tool
  - Challenge is to make them suitable for sensor platform

# Conclusion

---

- ❑ Making the sensor system robust should not be an after-thought
- ❑ If you need more details on
  - DADA → Shengchao Yu's poster
  - TARA → Jaewon Kang's talk in the afternoon
  - MARA → Badri Nath's talk in the afternoon

Questions ???