

802.11 Distributed Coordination Function: An ORBIT Case Study

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Outline

- ◆ Motivation and goals
- ◆ Experimental setup
- ◆ The first case study: adaptive PHY rate
- ◆ The second case study: fixed PHY rate
- ◆ Conclusions

Motivation

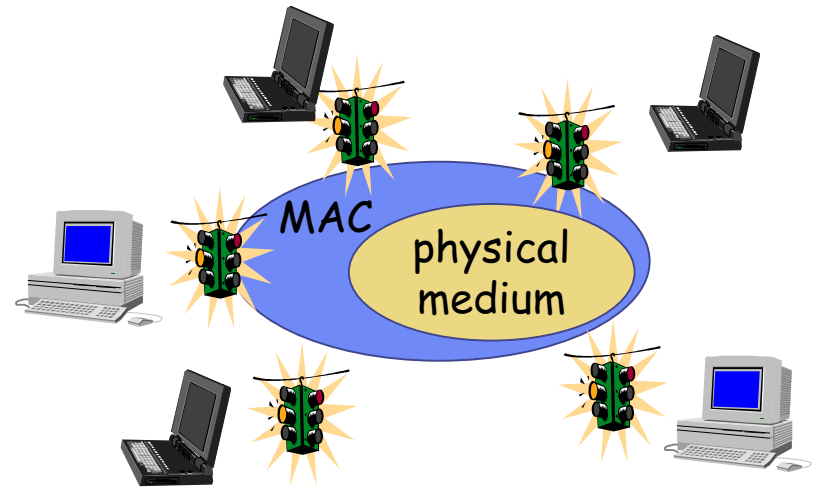
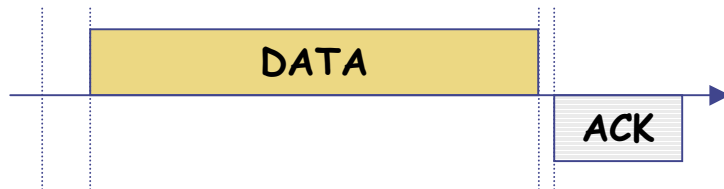
- ◆ The protocol is complex and hard to analyze analytically
- ◆ Analytical models [Bianchi'00, Cali et al '00] consider:
 - saturation - always a packet available for transmission
 - ideal channel - zero bit error probability
- ◆ Simulation studies trade model accuracy for processing effort
- ◆ Prior experimental work considers simple topologies and a small number of nodes [Kamerman '00, Anastasi et al '03 and '04, ...]

802.11 Distributed Coordination Function

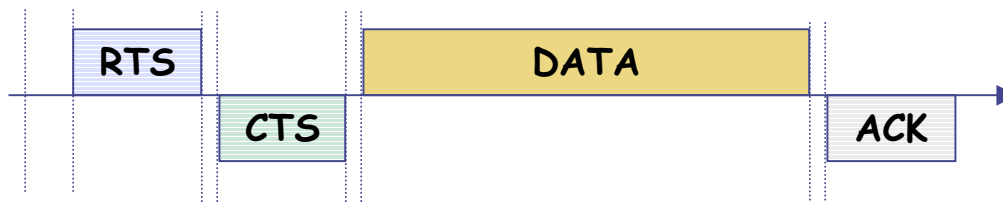
- ◆ MAC arbitrates access to the physical medium

- ◆ DCF is a random access scheme:

- basic access



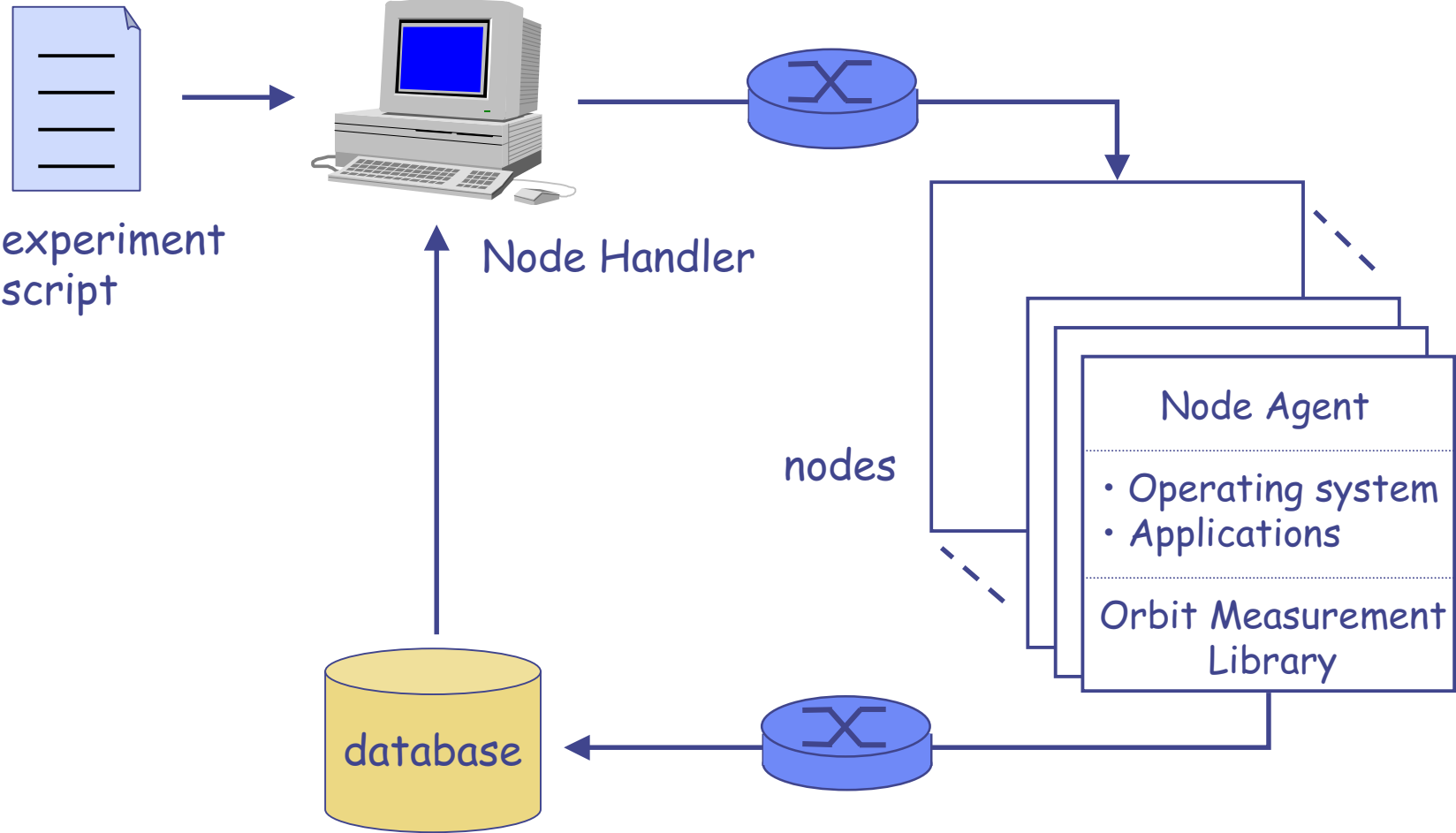
- RTS/CTS



Goals

- ◆ Test the network under different conditions:
 - PHY rate mechanism
 - ◆ adaptive rate
 - ◆ fixed rate
 - network topology
 - ◆ point-to-point
 - ◆ sink
 - network size
 - offered load
 - ◆ payload lengths
 - ◆ packet rates
- ◆ Compare results to the analytical model for saturation throughput from [Bianchi '00]

ORBIT testbed



Framework common to all experiments

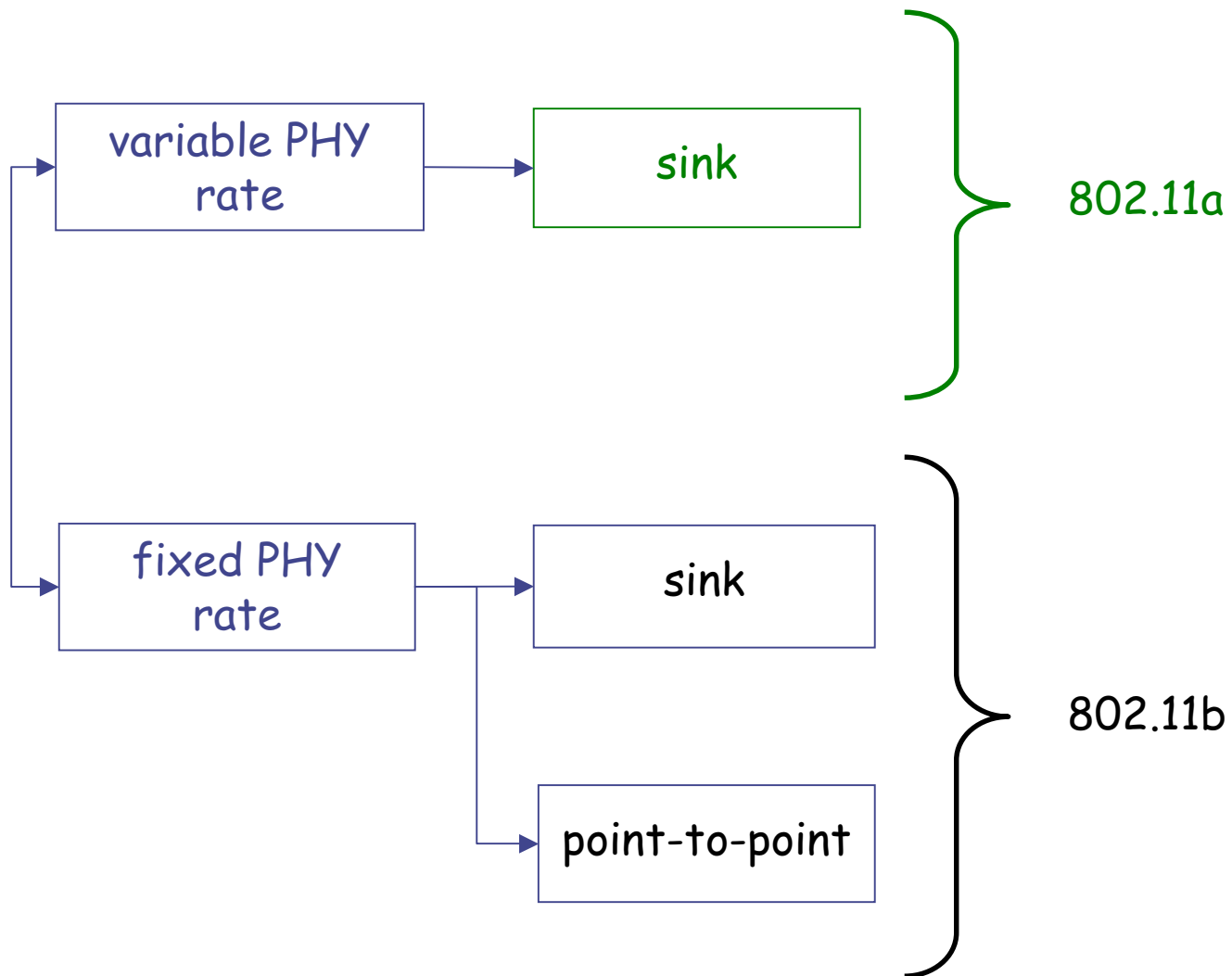
Aggregate performance on the transport layer:

- ◆ aggregate offered load (for all senders)
- ◆ aggregate throughput (for all receivers)
- ◆ average latency for successfully delivered packets

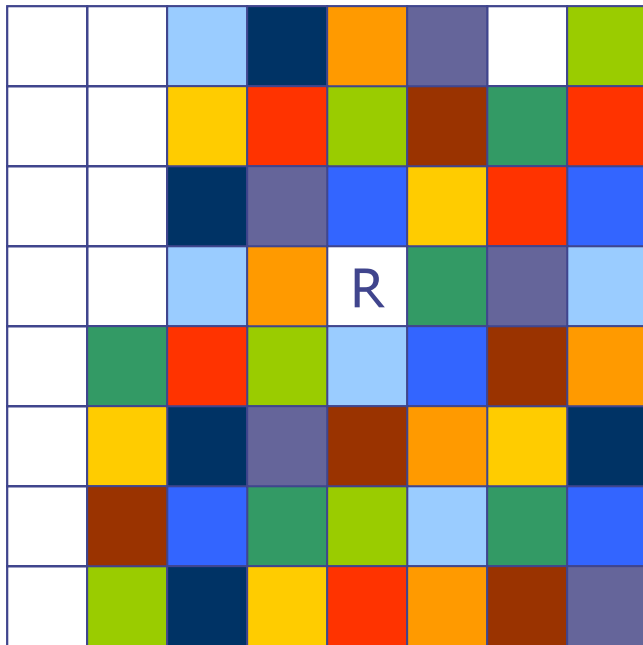
Common features:

- ◆ UDP CBR traffic
- ◆ high SNR due to the physical proximity
- ◆ fixed transmit power
- ◆ measurements collected every second
- ◆ settings fixed for 120 seconds

ORBIT case studies



First case study: Adaptive PHY rate

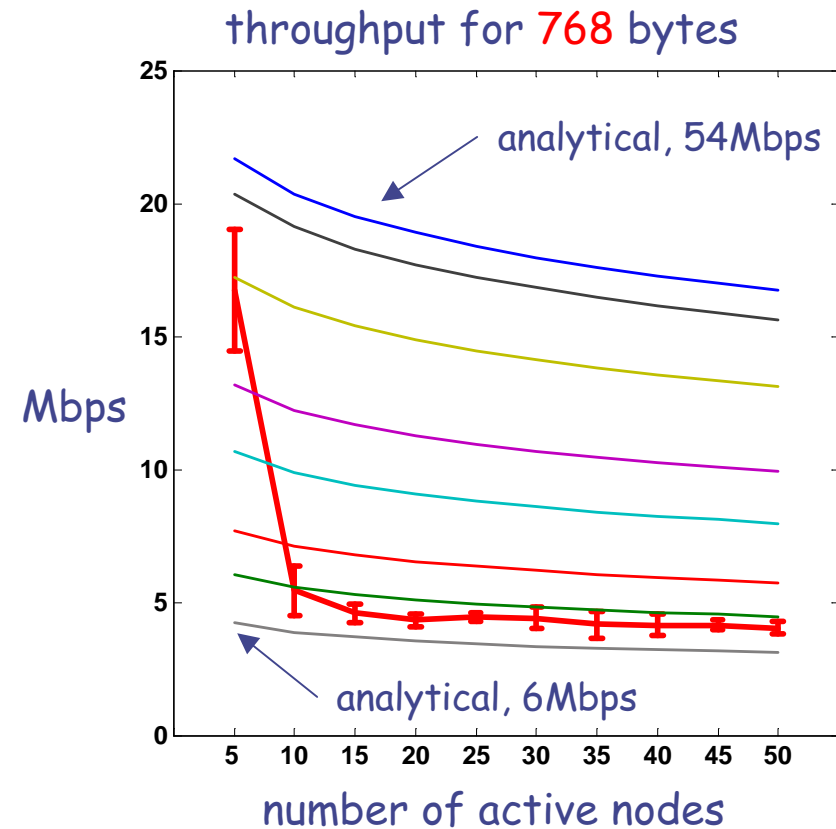
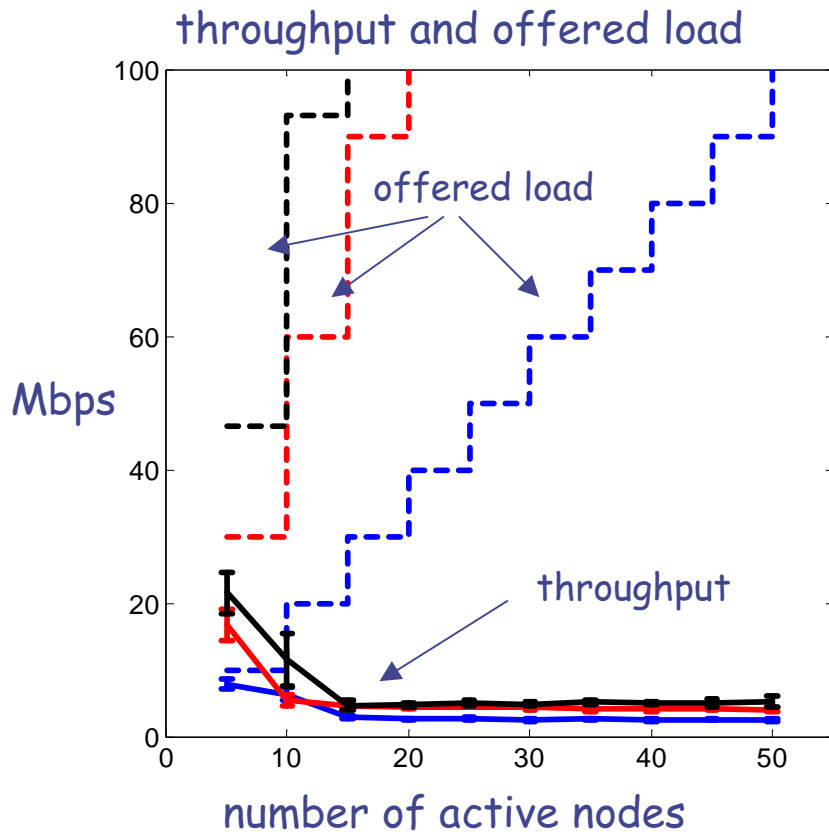


Saturated network:

- every 2 minutes add 5 senders
- total 50 senders

Throughput versus the number of nodes

802.11a basic access with payloads 256, 768, and 1280 bytes



Adaptive PHY rate algorithm

- ◆ Algorithm operation:
 - count missing acknowledgements
 - adjust the rate according to the number of failures
- ◆ Assumes low SNR as the source of packet loss
- ◆ Does not take into account DCF as a source of packet loss
- ◆ Conjecture:
PHY rate should be fixed in a congested network

The second case study: fixed PHY rate

topologies for the fixed PHY rate study

point-to-point

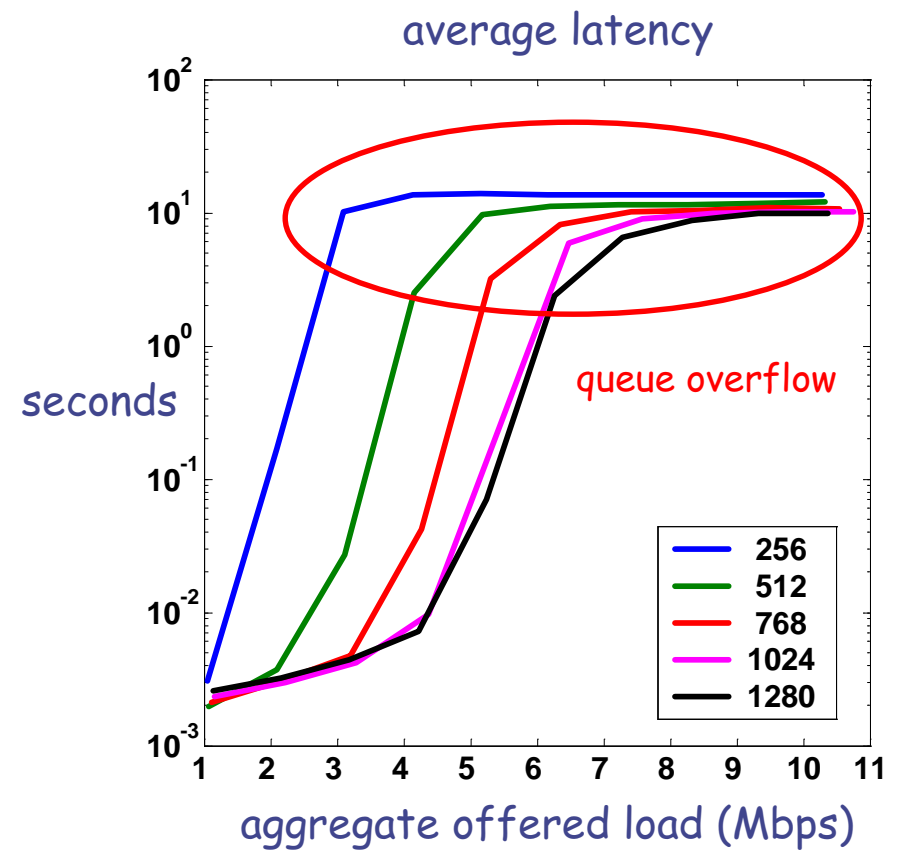
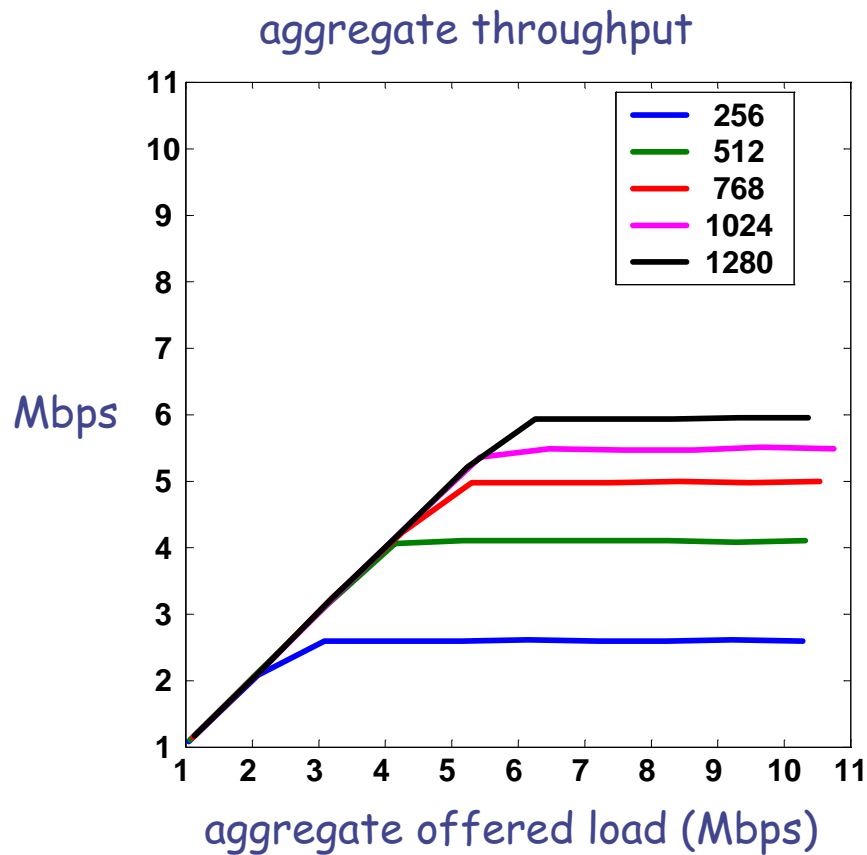
					2		2
			1		3		3
	1		4		4		
					5		5

sink

	S ₂				S ₇		S ₁₀
			S ₄		R		S ₉
	S ₁		S ₃		S ₆		
					S ₅		S ₈

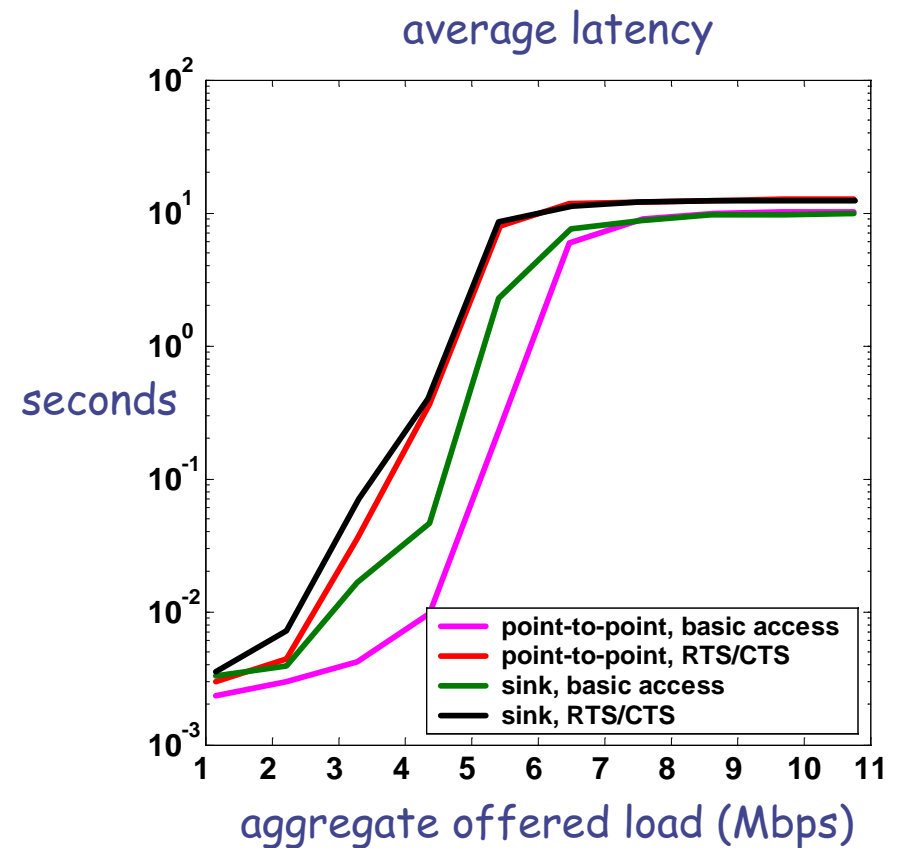
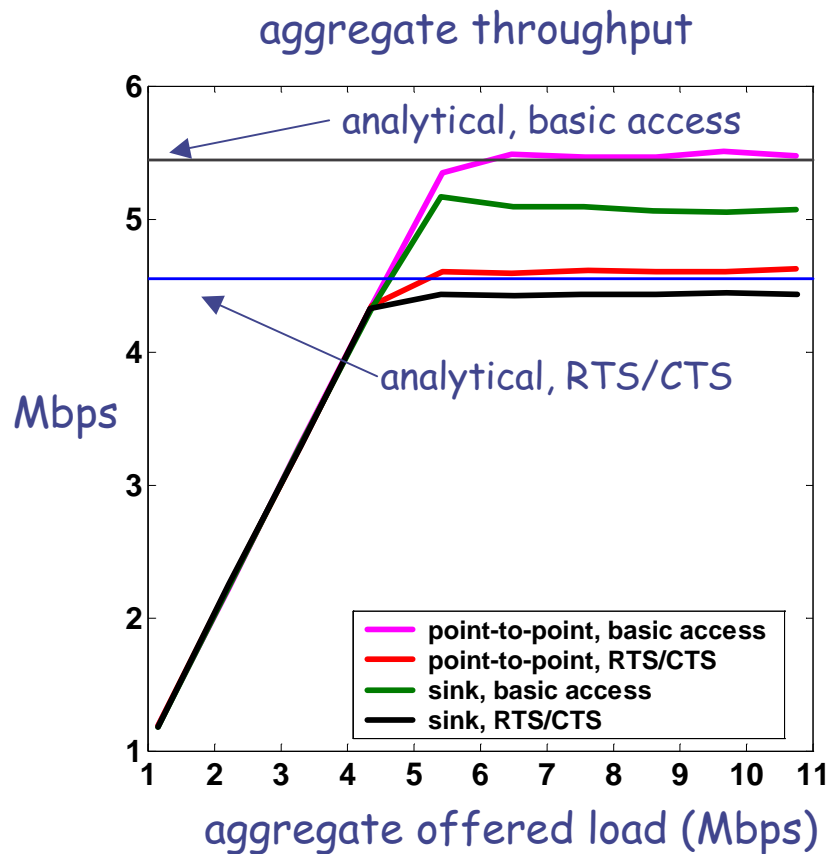
Throughput and latency versus offered load

802.11b basic access, rate 11Mbps, point-to-point



Throughput and latency: Basic access versus RTS/CTS

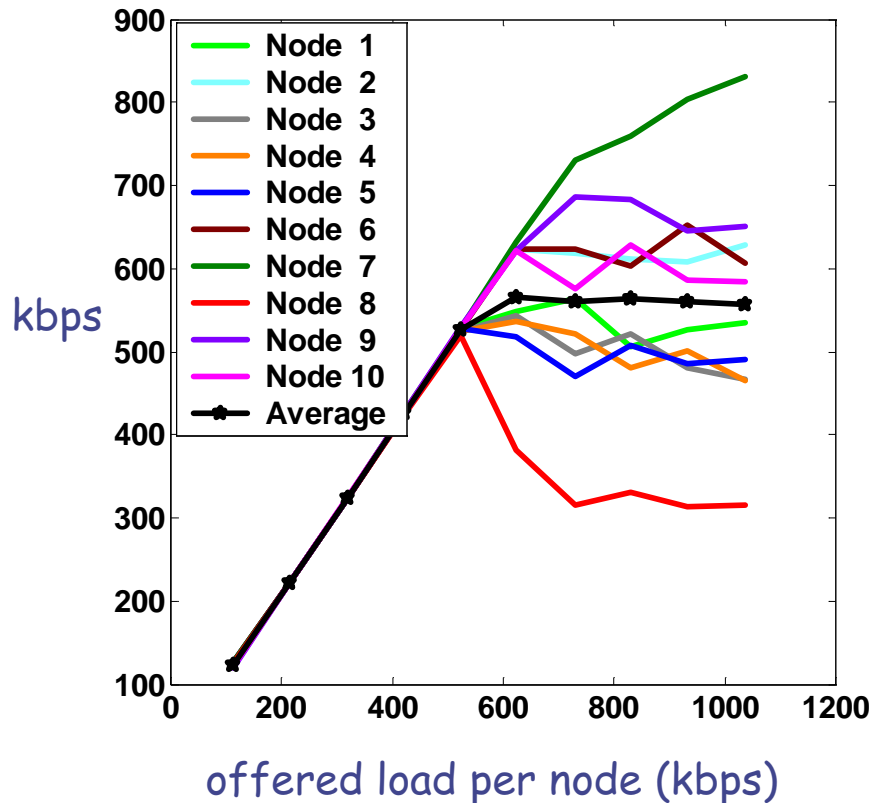
802.11b, payload 1024 bytes, rate 11Mbps



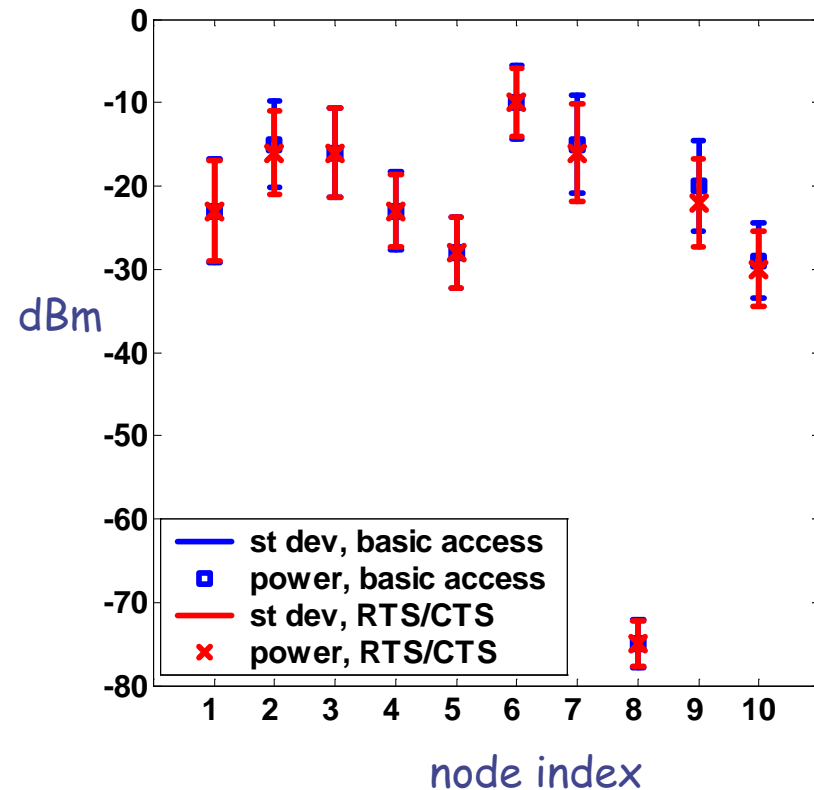
Throughput per node

- number of acknowledged packets multiplied with payload length in bits
- no fairness in saturation because of the capture effect

sink, basic access, payload 1280 bytes

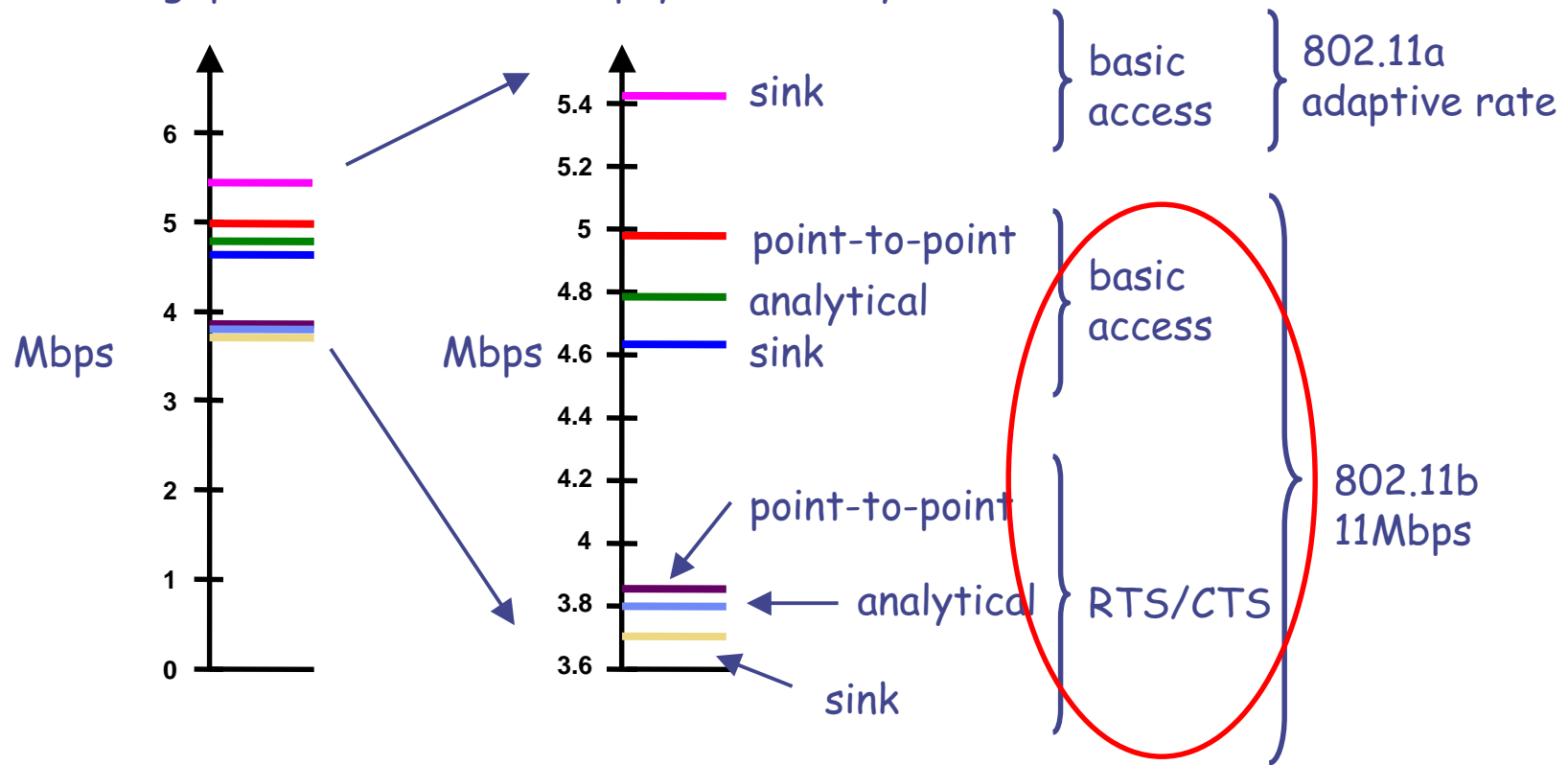


signal strength at the senders



Saturation throughput: The comparison of two case studies

throughput for 10 nodes and payload 768 bytes



basic access outperforms RTS/CTS

Conclusions

The study with adaptive PHY rate shows:

- ◆ adaptive PHY rate degrades the throughput in a congested network

The study with fixed PHY rate shows:

- ◆ for a relatively small network basic access outperforms RTS/CTS
- ◆ latency grows exponentially with an increase in the aggregate offered load
- ◆ in saturation nodes unequally contribute to the aggregate throughput
 - caused by the difference in signal strengths

Both studies show:

- ◆ [Bianchi '00] predicts well saturation throughput for networks with up to 50 nodes